On Process Rate Semantics

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Extract of the:
Open Lectures for PhD Students in Computer Science
Warsaw 2009-03-12...13

http://lucacardelli.name

Semantics of Collective Behavior

"Micromodels": Continuous Time Markov Chains

- The underlying semantics of stochastic π -calculus (and stochastic interacting automata). Well established in many ways.
 - Automata with rates on transitions.
- "The" correct semantics for chemistry, executable.
 - Gillespie stochastic simulation algorithm
- Lots of advantages
 - Compositional, compact, mechanistic, etc.
- But do not give a good sense of "collective" properties.
 - Yes one can do simulation.
 - Yes one can do program analysis.
 - Yes one can perhaps do modelchecking.
 - But somewhat lacking in "analytical properties" and "predictive power".

"Macromodels": Ordinary Differential Equations

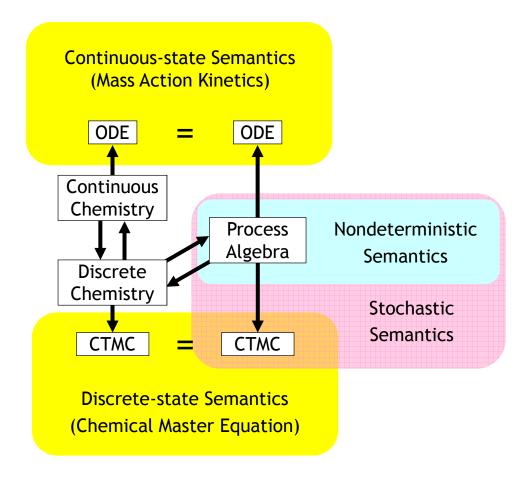
- The classical semantics of collective behavior.
 - E.g. kinetic theory of gasses.
 - They always ask: "How does you automata model relate to the 75 ODE models in the literature?"
- Going from processes/automata to ODEs <u>directly</u>:
 - o *In principle*: just write down the Rate Equation:
 - Let [S] be the "number of processes in state S" as a function of time.
 - Define for each state S:

```
d[S]/dt = (rate of change of the number of processes in state S)
    Cumulative rate of transitions from any state S' to state S, times [S'],
    minus cumulative rate of transitions from S to any state S", times [S].
```

- Fairly intuitive (rate = inflow minus outflow)
- Going to ODEs <u>indirectly</u> through chemistry
 - If we first convert processes to chemical reactions, then we can convert to ODEs by standard means!



The Two Semantic Sides of Chemistry

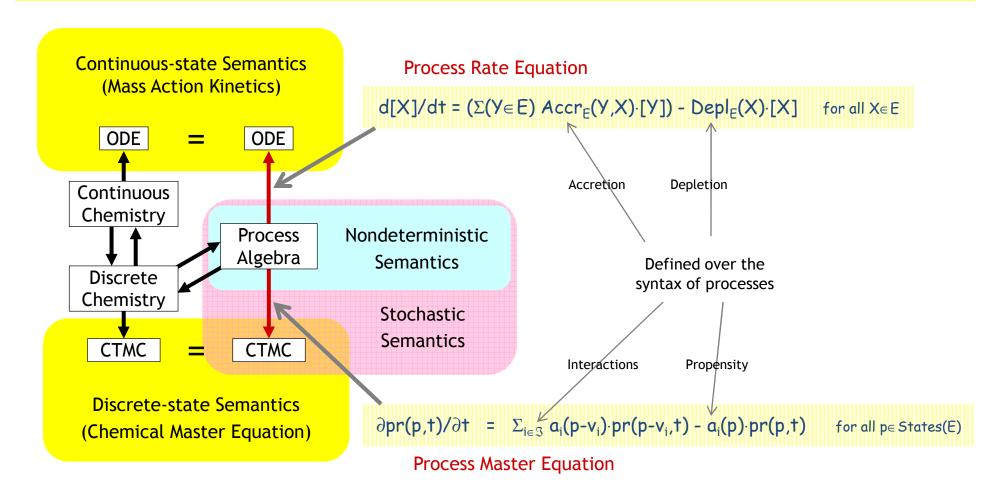


These diagrams commute via appropriate maps.

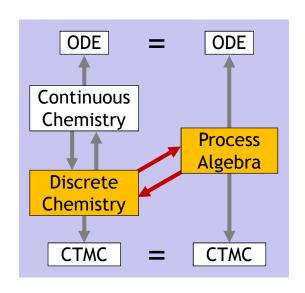
L. Cardelli: "On Process Rate Semantics" (TCS)

L. Cardelli: "A Process Algebra Master Equation" (QEST'07)

Quantitative Process Semantics



Stochastic Processes & Discrete Chemistry



Chemical Reactions (FSRN)

$$A \longrightarrow^{r} B_1 + ... + B_n (n \ge 0)$$

$$A_1 + A_2 \rightarrow^r B_1 + ... + B_n$$
 (n≥0) Hetero Reaction

$$A + A \longrightarrow^r B_1 + ... + B_n (n \ge 0)$$

Unary Reaction

Homeo Reaction

d[A]/dt = -r[A]

Exponential Decay

$$d[A_i]/dt = -r[A_1][A_2]$$

Mass Action Law

$$d[A]/dt = -2r[A]^2$$

Mass Action Law

(assuming $A \neq B_i \neq A_i$ for all i,j)

No other reactions!

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The chemical Langevin equation

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Genuinely trimolecular reactions do not physically occur in dilute fluids with any appreciable frequency. Apparently trimolecular reactions in a fluid are usually the combined result of two bimolecular reactions and one monomolecular reaction, and involve an additional short-lived species.

Chapter IV: Chemical Kinetics [David A. Reckhow, CEE 572 Course]

... reactions may be either elementary or nonelementary. Elementary reactions are those reactions that occur exactly as they are written, without any intermediate steps. These reactions almost always involve just one or two reactants. ... Non-elementary reactions involve a series of two or more elementary reactions. Many complex environmental reactions are nonelementary. In general, reactions with an overall reaction order greater than two, or reactions with some non-integer reaction order are non-elementary.

THE COLLISION THEORY OF REACTION RATES www.chemguide.co.uk

The chances of all this happening if your reaction needed a collision involving more than 2 particles are remote. All three (or more) particles would have to arrive at exactly the same point in space at the same time, with everything lined up exactly right, and having enough energy to react. That's not likely to happen very often!

Trimolecular reactions:

$$A + B + C \rightarrow^r D$$

the measured "r" is an (imperfect) aggregate of e.g.:

$$A + B \leftrightarrow AB$$

$$AB + C \rightarrow D$$

Enzymatic reactions:

$$S \xrightarrow{E} P$$

the "r" is given by Michaelis-Menten (approximated steady-state) laws:

$$E + S \leftrightarrow ES$$

$$ES \rightarrow P + E$$

Chemical Ground Form (CGF)

E ::= 0 : X=M, E Reagents

 $M := 0 : \pi; P \oplus M$ Molecules

P::= 0 : X | P Solutions

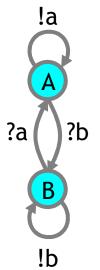
 $\pi ::= \tau_{(r)} : ?a_{(r)} : !a_{(r)}$ Actions (delay, input, output)

CGF ::= E,P Reagents plus Initial Conditions

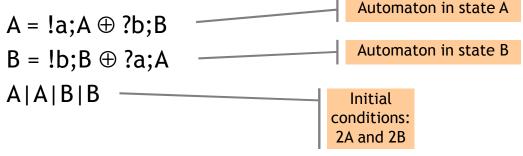
A stochastic subset of CCS (no values, no restriction)

(To translate chemistry to processes we need a bit more than interacting automata: we may have "+" on the right of \rightarrow , that is we may need "|" after π .)

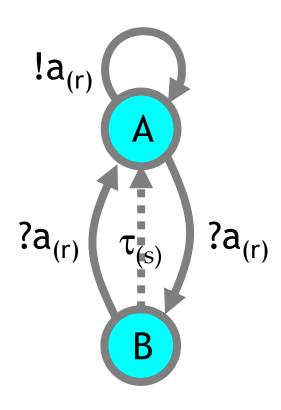
 \oplus is stochastic choice (vs. + for chemical reactions) 0 is the null solution (P|0 = 0|P = P) and null molecule (M \oplus 0 = 0 \oplus M = M) Each X in E is a distinct *species* Each name a is assigned a fixed rate r: $a_{(r)}$



Ex: Interacting Automata (= finite-control CGFs: they use "|" only in initial conditions):

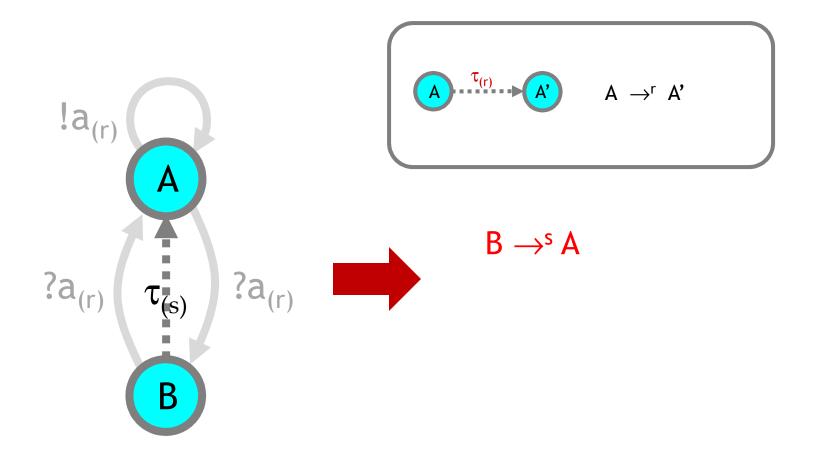


From CGF to Chemistry



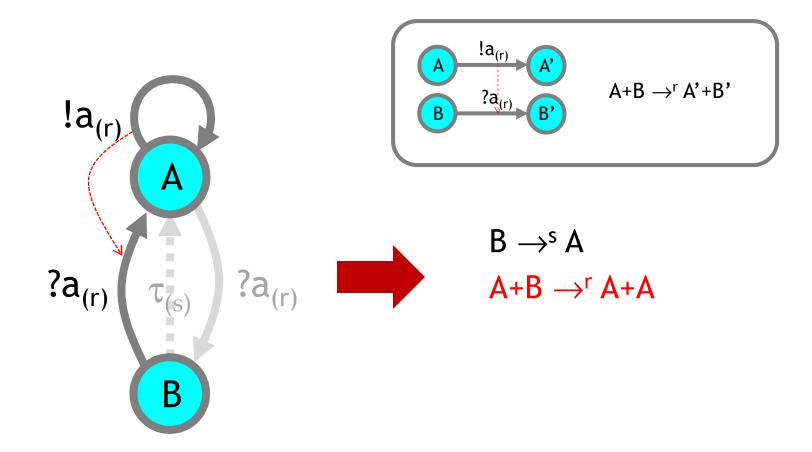
$$A = !a_{(r)};A \oplus ?a_{(r)};B$$

$$B = ?a_{(r)};A \oplus \tau_{(s)};A$$



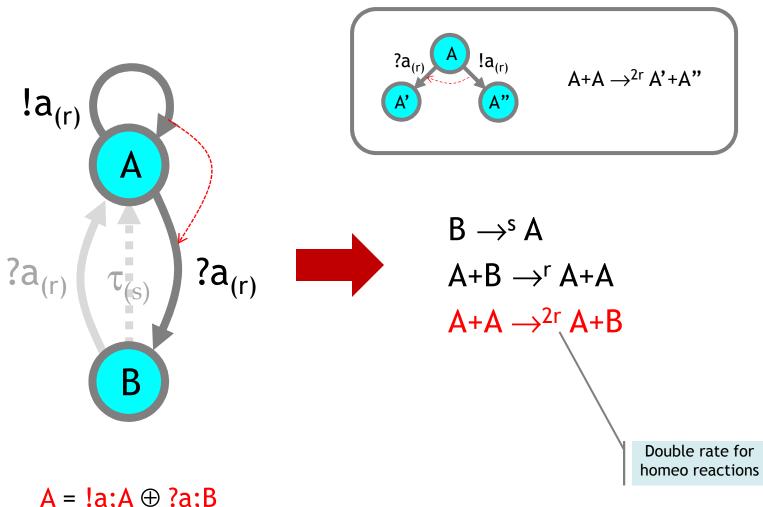
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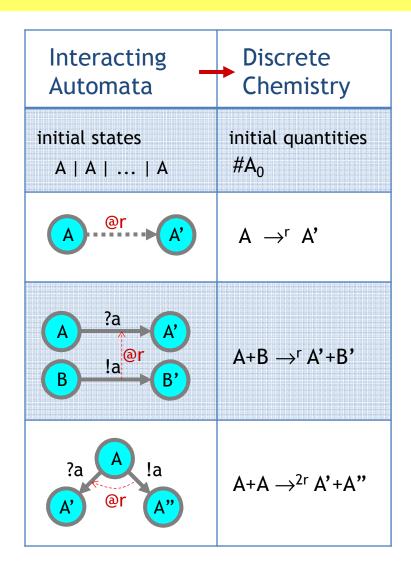
$$A = !a; A \oplus ?a; B$$

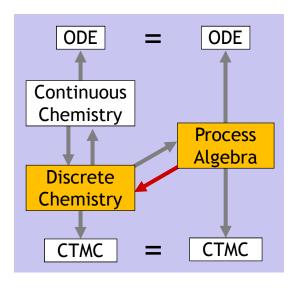
$$B = ?a;A \oplus \tau_{(s)};A$$



 $A = !a;A \oplus ?a;B$

 $B = ?a;A \oplus \tau_{(s)};A$





From CGF to Chemistry: Ch(E)

E::= 0 : X=M, E Reagents

 $M := 0 : \pi; P \oplus M$ Molecules

P := 0 : X | P Solutions

 $\pi ::= \tau_{(r)} : ?a_{(r)} : !a_{(r)}$ Interactions (delay, input, output)

CGF ::= E,P Reagents plus Initial Conditions

E.X.i == the i-th Å-summand of the molecule M associated with the X reagent of E

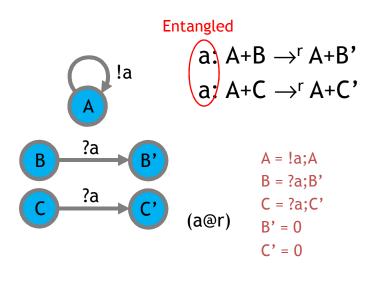
Chemical reactions for E,P: (N.B.: <...> are reaction tags to obtain multiplicity of reactions, and P is P with all the | changed to +)

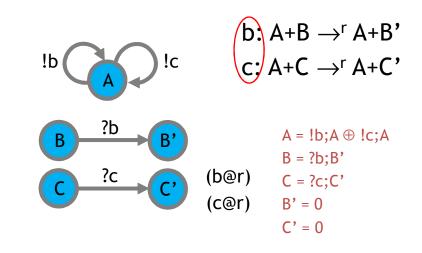
```
Ch(E) :=  \{(\langle X.i \rangle: X \to^{r} P) \text{ s. t. E.X.} i = \tau_{(r)}; P\} \cup \\ \{(\langle X.i, Y.j \rangle: X + Y \to^{r} P + Q) \text{ s. t. } X \neq Y, \text{ E.X.} i = ?a_{(r)}; P, \text{ E.Y.} j = !a_{(r)}; Q\} \cup \\ \{(\langle X.i, X.j \rangle: X + X \to^{2r} P + Q) \text{ s. t. E.X.} i = ?a_{(r)}; P, \text{ E.X.} j = !a_{(r)}; Q\}
```

Initial conditions for P:

Ch(P) := P

Entangled vs Detangled





Entangled: Two reactions on one channel

Detangled: Two reactions on two separate channels

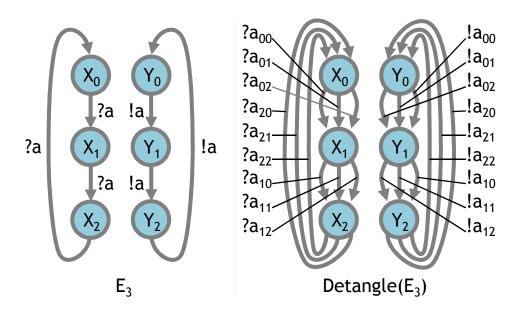
We need a semantics of automata that identifies automata that have the "same chemistry".

No traditional process algebra equivalence is like this!

Entangled automata lead to more compact models than in chemistry.

Detangled automata are in simple correspondence with chemistry.

Entangled vs detangled



(closely related to $Pi(Ch(E_3))$)

Chemical Parametric Form (CPF)

E := 0 : X(p)=M, E

 $M ::= 0 : \pi; P \oplus M$

P := 0 : X(p) | P

 $\pi ::= \tau_{(r)} : ?a_{(r)}(\mathbf{p}) : !a_{(r)}(\mathbf{p})$

CPF::= **E**,**P**

Not bounded-state systems.

Not finite-control systems.

But still finite-species systems.

A translation from CPF to CGF exists (expanding all possible instantiation of parameters from the initial conditions)

An incremental translation algorithm exists (expanding on demand from initial conditions)

Reagents

Molecules

Solutions

Actions

with initial conditions

⊕ is stochastic choice (vs. + for chemical reactions)

0 is the null solution (P|0 = 0|P = P)and null molecule $(M \oplus 0 = 0 \oplus M = M)$

Each X in E is a distinct species

p are vectors of names

p are vectors of distinct names when in binding position

Each free name a in E is assigned a fixed rate r: $a_{(r)}$

Example:

 $Neg(a,b) = ?a; Inh(a,b) \oplus \tau_e; (Tr(b) | Neg(a,b))$

 $Inh(a,b) = \tau_h; Neg(a,b)$

 $Tr(b) = !b; Tr(b) \oplus \tau_d; 0$

Neg(x,x)

CPF to CGF: Handling Parameters

Consider first the CPF subset with no communication (pure ?a, !a).

Grounding (replace parameters with constants)

where X/p is a name in bijection with <X,p> (each X/p is seen as a separate *species*)

$$/(\pi_1; P_1 \oplus ... \oplus \pi_n; P_n) =_{def} \pi_1; /(P_1) \oplus ... \oplus \pi_n; /(P_n)$$
 $/(X_1(p_1) | ... | X_n(p_n)) =_{def} X_1/p_1 | ... | X_n/p_n$

E ::=
$$X_1(\mathbf{p}_1) = M_1$$
, ..., $X_n(\mathbf{p}_n) = M_n$
M ::= $\pi_1; P_1 \oplus ... \oplus \pi_n; P_n$
P ::= $X_1(\mathbf{p}_1) \mid ... \mid X_n(\mathbf{p}_n)$
 $\pi ::= \tau_r$?a !a

Let N be the set of free names occurring in E.

E_G is the **Parametric Explosion** of E (still a finite species system) computed by replacing parameters with all combinations of free names in E

$$E_G := \{(X/\mathbf{q} = /(M\{\mathbf{p}\leftarrow\mathbf{q}\})) \text{ s.t. } (X(\mathbf{p}) = M) \in E \text{ and } \mathbf{q} \in N^{\#p}\}$$

$$P_G := /P \qquad \text{(simply ground the given initial conditions once)}$$

 E_G is a CGF! To obtain the chemical reactions $Ch_P(E)$, just compute $Ch_G(E_G)$

$$Ch_P(E) = Ch_G(E_G)$$

CPF to CGF: Handling Communication

Grounding (replace parameters with constants)

just one main change: now also convert each input parameter into a ground choice of all possible inputs

N is the set of free names in E,P

#p is the length of p

n/p is a name in bijection with <n,p>

X/p is a name in bijection with <X,p>

(each X/p is seen as a separate species)

$$/_{N}(\tau_{r};P) = \tau_{r}; /_{N}(P)$$

$$/_{N}(!a_{(r)}(p);P) = !a/p_{(r)}; /_{N}(P)$$

$$/_{N}(?a_{(r)}(p);P) = \oplus (q \in N^{\#p}) \text{ of } ?a/q_{(r)}; /_{N}(P\{p \leftarrow q\})$$

$$/_{N}(\pi_{1};P_{1} \oplus ... \oplus \pi_{n};P_{n}) = /_{N}(\pi_{1};P_{1}) \oplus ... \oplus /_{N}(\pi_{n};P_{n})$$

$$/_{N}(X_{1}(p_{1}) \mid ... \mid X_{n}(p_{n})) = X_{1}/p_{1} \mid ... \mid X_{n}/p_{n}$$

E ::=
$$X_1(p_1)=M_1$$
, ..., $X_n(p_n)=M_n$
M ::= $\pi_1; P_1 \oplus ... \oplus \pi_n; P_n$
P ::= $X_1(p_1) \mid ... \mid X_n(p_n)$
 $\pi ::= \tau_r$?a(p) !a(p)

E_G is again the **Parametric Explosion** of E

$$\begin{split} E_G &:= \{ (X/\mathbf{q} = /_N(M\{p \leftarrow \mathbf{q}\})) \text{ s.t. } (X(\mathbf{p}) = M) \in E \text{ and } \mathbf{q} \in N^{\#p} \} \\ P_G &:= /_N(P) \end{split} \qquad \text{(simply ground the given initial conditions once)} \end{split}$$

$$Ch(E) = Ch_G(E_G)$$
 E_G is a again a CGF!

CPF to CGF Translation. Ex: Neg(x,x)

```
E =
    Neg(a,b) = ?a; Inh(a,b) \oplus \tau_e; (Tr(b) \mid Neg(a,b))
    Inh(a,b) = \tau_h; Neg(a,b)
    Tr(b) = !b; Tr(b) \oplus \tau_d; 0
    Neg(x,x)
---- initialization -----
E_c := \{ \text{Neg/x}, x = ?x; \text{Inh/x}, x \oplus \tau_e; (\text{Tr/x} \mid \text{Neg/x}, x) \}
---- iteration 1 -----
C := \{ \text{Neg/x}, x \rightarrow^{\text{e}} \text{Tr/x} + \text{Neg/x}, x \}
E_c := \{ \text{Neg/x,x} = ?x; \text{Inh/x,x} \oplus \tau_e; (\text{Tr/x} \mid \text{Neg/x,x}) \}
         Tr/x = !x; Tr/x \oplus \tau_d; 0
---- iteration 2 -----
C := \{ \text{Neg/x}, x \rightarrow^{\text{e}} \text{Tr/x} + \text{Neg/x}, x \}
         Tr/x \rightarrow d 0
         Tr/x + Neg/x, x \rightarrow^{\rho(x)} Tr/x + Inh/x, x 
E_c := \{ \text{Neg/x,x} = ?x; \text{Inh/x,x} \oplus \tau_e; (\text{Tr/x} \mid \text{Neg/x,x}) \}
         Tr/x = !x; Tr/x \oplus \tau_d; 0
         lnh/x,x = \tau_h; Neg/x,x
```

```
:---- iteration 3 -----
C := \{ Neg/x, x \rightarrow^e Tr/x + Neg/x, x \\ Tr/x \rightarrow^d 0 \\ Tr/x + Neg/x, x \rightarrow^{\rho(x)} Tr/x + lnh/x, x \\ lnh/x, x \rightarrow^h Neg/x, x \}
E_c := \text{ no change}
:---- termination -----
Neg/x, x \rightarrow^e Tr/x + Neg/x, x \\ Tr/x \rightarrow^d 0 \\ Tr/x + Neg/x, x \rightarrow^{\rho(x)} Tr/x + lnh/x, x \\ lnh/x, x \rightarrow^h Neg/x, x \\ Neg/x, x
```

From Chemistry to CGF

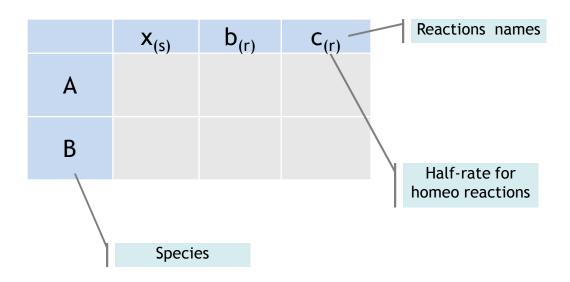
From Chemistry to CGF (by example)

 $x: B \rightarrow^s A$

b: $A+B \rightarrow^r A+A$

c: $A+A \rightarrow^{2r} A+B$

Unique reaction names



From Chemistry to CGF (by example)

$$x: B \rightarrow S A$$

b:
$$A+B \rightarrow^r A+A$$

c:
$$A+A \rightarrow^{2r} A+B$$

	X (s)	$b_{(r)}$	c _(r)
Α			
В	τ;Α		

1: Fill the matrix by columns:

Degradation reaction v_i : $X \rightarrow k_i P_i$ add τ ; P_i to $\langle X, v_{ij} \rangle$.

$$x: B \rightarrow^s A$$

b:
$$A+B \rightarrow^r A+A$$

c:
$$A+A \rightarrow^{2r} A+B$$

	X (s)	b _(r)	C _(r)
Α		?;A A	
В	τ;Α	!;0	

1: Fill the matrix by columns:

```
Degradation reaction v_i: X \rightarrow k_i P_i
add \tau;P_i to < X, v_{ij} >.
Hetero reaction v_i: X+Y \rightarrow k_i P_i
add ?;P_i to < X, v_i > and !;0 to < Y, v_i >
```

$$x: B \rightarrow^s A$$

b:
$$A+B \rightarrow^r A+A$$

c:
$$A+A \rightarrow 2r A+B$$

	X (s)	$b_{(r)}$	C (r)
Α		?;A A	?;A B !;0
В	τ;Α	!;0	

1: Fill the matrix by columns:

```
Degradation reaction v_i: X \rightarrow k_i P_i
add \tau; P_i to \langle X, v_{ij} \rangle.
Hetero reaction v_i: X+Y \rightarrow k_i P_i
add ?; P_i to \langle X, v_i \rangle and !; 0 to \langle Y, v_i \rangle
Homeo reaction v_i: X+X \rightarrow k_i P_i
add ?; P_i and !; 0 to \langle X, v_i \rangle
```

$$x: B \rightarrow^s A$$

b:
$$A+B \rightarrow^r A+A$$

c:
$$A+A \rightarrow 2^r A+B$$

$x_{(s)}$ $b_{(r)}$ $c_{(r)}$?;A|B ?;A|A $\tau;A$!;0

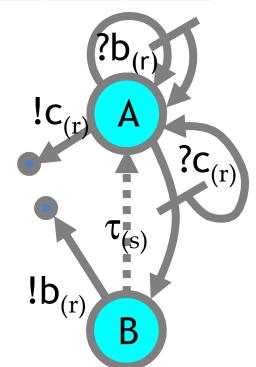
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Hetero reaction v_i : $X+Y \rightarrow k_i P_i$
add ?; P_i to $< X, v_i >$ and !; 0 to $< Y, v_i >$
Homeo reaction v_i : $X+X \rightarrow k_i P_i$
add ?; P_i and !; 0 to $< X, v_i >$

2: Read the result by rows:

$$A = ?b_{(r)}; (A|A) \oplus ?c_{(r)}; (A|B) \oplus !c_{(r)}; 0$$

 $B = \tau_{(s)}; A \oplus !b_{(r)}; 0$



$$x: B \rightarrow^s A$$

b:
$$A+B \rightarrow^r A+A$$

c:
$$A+A \rightarrow^{2r} A+B$$

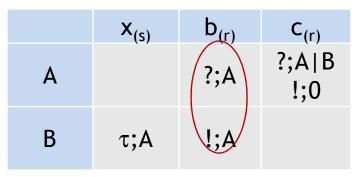
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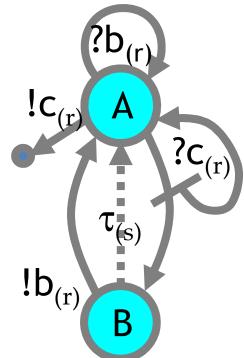
Degradation reaction
$$v_i$$
: $X \rightarrow k_i P_i$
add τ ; P_i to $< X, v_{ij} >$.
Hetero reaction v_i : $X+Y \rightarrow k_i P_i$
add $?$; P_i to $< X, v_i >$ and $!$; 0 to $< Y, v_i >$
Homeo reaction v_i : $X+X \rightarrow k_i P_i$
add $?$; P_i and $!$; 0 to $< X, v_i >$

2: Read the result by rows:

$$A = ?b_{(r)}; A \oplus ?c_{(r)}; (A | B) \oplus !c_{(r)}; 0$$

$$B = \tau_{(s)}; A \oplus !b_{(r)}; A$$





$$x: B \rightarrow^s A$$

b:
$$A+B \rightarrow^r A+A$$

c:
$$A+A \rightarrow^{2r} A+B$$

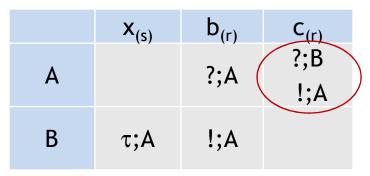
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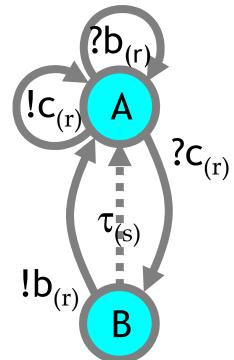
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Hetero reaction v_i : $X+Y \rightarrow k_i P_i$
add $?$; P_i to $< X, v_i >$ and $!$; 0 to $< Y, v_i >$
Homeo reaction v_i : $X+X \rightarrow k_i P_i$
add $?$; P_i and $!$; 0 to $< X, v_i >$

2: Read the result by rows:

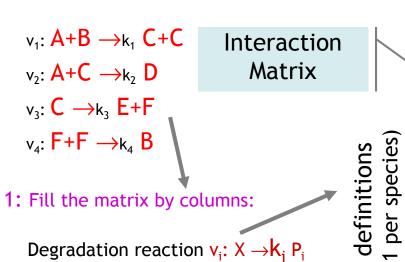
$$A = ?b_{(r)}; A \oplus ?c_{(r)}; B \oplus !c_{(r)}; A$$

$$B = \tau_{(s)}; A \oplus !b_{(r)}; A$$





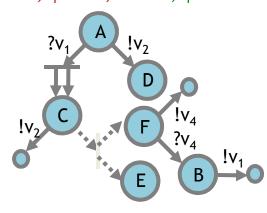
From Chemistry to Automata (by example)



Degradation reaction $v_i: X \rightarrow k_i P_i$ add τ ; P_i to $\langle X, V_{ij} \rangle$.

Hetero reaction $v_i: X+Y \rightarrow k_i P_i$ add ?; P_i to $\langle X, v_i \rangle$ and !;0 to $\langle Y, v_i \rangle$

Homeo reaction $v_i: X+X \rightarrow k_i P_i$ add ?; P_i and !;0 to $\langle X, v_i \rangle$



channels and rates (1 per reaction)

	V _{1(k1)}	$V_{2(k2)}$	$V_{3(k3)}$	V _{4(k4/2)}
Α	?;(C C)	?;D		
В	!;0			
С		!;0	τ;(E F)	
D				
Ε				
F				?;B !;0

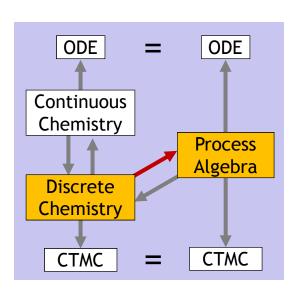
2: Read the result by rows:

 $A = ?v_{1(k1)}; (C | C) \oplus ?v_{2(k2)}; D$

 $B = !v_{1(k1)};0$

 $C = !v_{2(k2)};0 \oplus \tau_{k3};(E|F)$

 $F = {}^{?}V_{4(k4/2)}; B \oplus {}^{!}V_{4(k4/2)}; 0$



Half-rate for

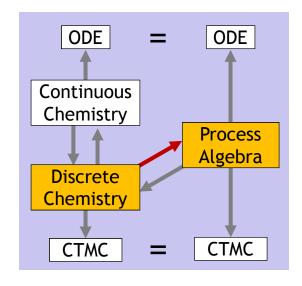
homeo reactions

From Chemistry to CGF: Pi(C)

v:
$$X \rightarrow^r Y_1 + ... + Y_n + 0$$
 Unary Reaction
v: $X_1 + X_2 \rightarrow^r Y_1 + ... + Y_n + 0$ Binary Reaction

From uniquely-labeled (v:) chemical reactions C to a CGF Pi(C):

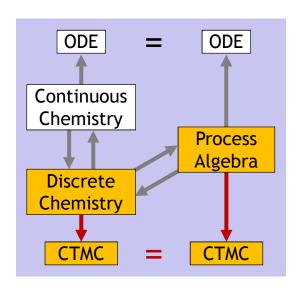
```
\begin{array}{ll} \text{Pi}(C) = \{(X = \oplus((v: X \rightarrow^k P) \in C) \ \textit{of} \ (\tau_{(k)}; P) & \oplus \\ & \oplus((v: X + Y \rightarrow^k P) \in C \ \text{and} \ Y \neq X) \ \textit{of} \ (?v_{(k)}; P) & \oplus \\ & \oplus((v: Y + X \rightarrow^k P) \in C \ \text{and} \ Y \neq X) \ \textit{of} \ (!v_{(k)}; 0) & \oplus \\ & \oplus((v: X + X \rightarrow^k P) \in C) \ \textit{of} \ (?v_{(k/2)}; P \oplus !v_{(k/2)}; 0) &) \\ & \textit{s.t.} \ X \ \text{is a species in } C\} \end{array}
```



Some Syntactic Properties

- C and Ch(Pi(C)) have the same reactions
 - o (and their reaction labels are in bijection)
- Def: E is detangled if each channel appears once as ?a and once as !a.
- If C is a system of chemical reactions then Pi(C) is detangled.
 - (hence chemical reactions embed into a subclass of CGFs)
- Hence for any E, we have that Pi(Ch(E)) is detangled.
 - (E and Pi(Ch(E)) are "equivalent" CGFs, but that has to be shown later)
- Def: E,P is automata form if "|" occurs only (other than "|0") in P.
- Def: Detangle(E) is defined from Pi(Ch(E)) by replacing any occurrence pairs $a_{(r)}$; (X|Y|0) and $a_{(r)}$; (X|0) and $a_{(r)}$; (X|0).
- If E is in automata form then Detangle(E) is (detangled and) in automata form
 - (but Pi(Ch(E)) may not be)

Discrete-State Semantics



Discrete Semantics of Reactions

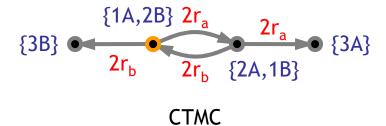
Syntax:

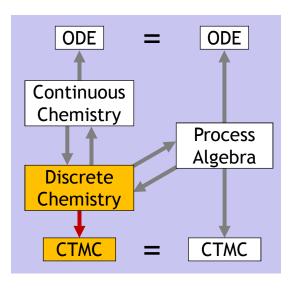
$$A+B \rightarrow^r A+A$$

$$A+B \rightarrow^r B+B$$

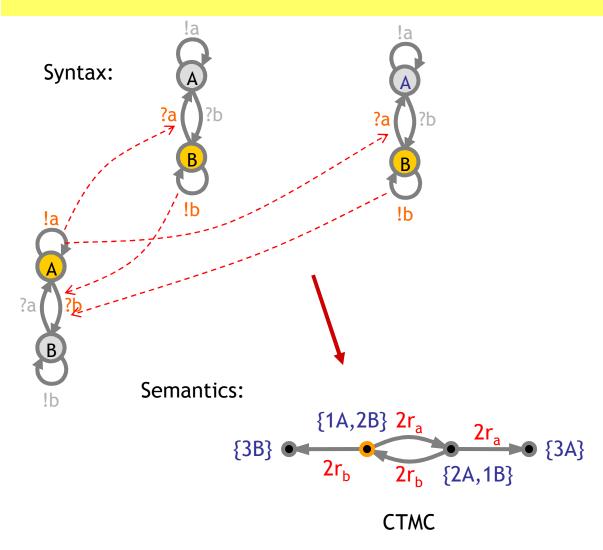


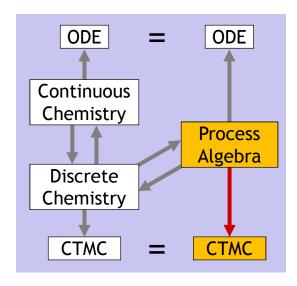
Semantics:





Discrete Semantics of Reagents

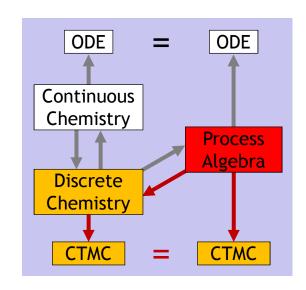


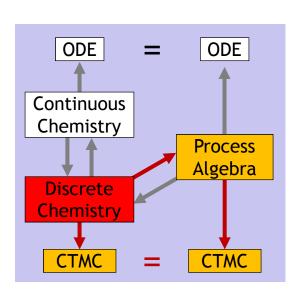


Discrete State Equivalence

Def:

 is equivalent CTMC's (isomorphic graphs with same rates).





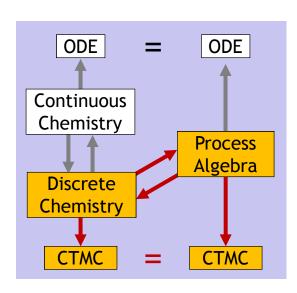
- For each E in automata form there is an an E'
 \(\times \) E that is detangled and in automata form (E' = Detangle(E)).

Interacting Automata = Discrete Chemistry

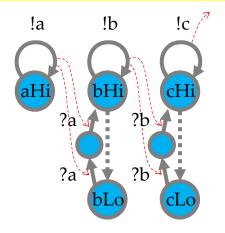
This is enough to establish that the process algebra is really faithful to the chemistry.

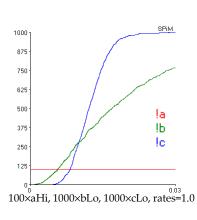
But CTMC are not the "ultimate semantics" because there are still questions of when two different CTMCs are actually equivalent (e.g. "lumping").

The "ultimate semantics" of chemistry is the *Chemical Master Equation* (derivable from the Chapman-Kolmogorov equation of the CTMC).

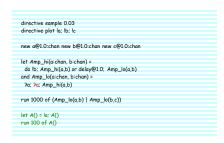


Exercise 1

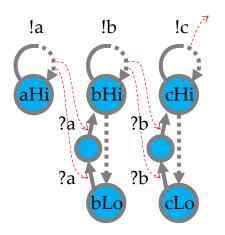


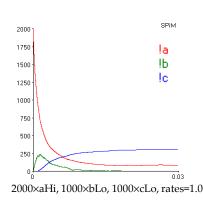


Second-Oder Regime cascade: a signal amplifier (MAPK) aHi > 0 \Rightarrow cHi = max



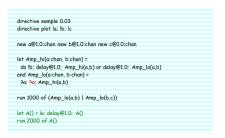
Write these automata in CGF and translate them to chemical reactions.



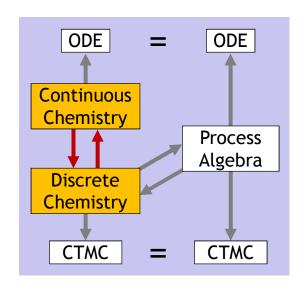


Zero-Oder Regime cascade: a signal divider!

 $aHi = max \Rightarrow cHi = 1/3 max$



Discrete vs Continuous Chemistry



The "Type System" of Chemistry

The International System of Units (SI) defines the following physical units, with related derived units and constants; note that amount of substance is a base unit in SI, like length and time:

```
mol (a base unit)mole, unit of amount of substancem (a base unit)meter, unit of lengths (a base unit)second, unit of timeL = 0.001 \cdot m^3liter (volume)M = mol \cdot L^{-1}molarity (concentration of substance)N_A : mol^{-1} \cong 6.022 \times 10^{23}Avogadro's number (number of particles per amount of substance)
```

For a substance X:mol, we write [X]:M for the concentration of X, and $[X]^{\bullet}:M\cdot s^{-1}$ for the time derivative of the concentration.

A continuous chemical system (C,V) is a system of chemical reactions C plus a vector of initial concentrations V_x : M, one for each species X.

The rates of unary reactions have dimension s⁻¹.

The rates of binary reactions have dimension M⁻¹s⁻¹.

(because in both cases the rhs of an ODE should have dimension M·s⁻¹).

Relating Concentration to Number of Molecules

For a given volume of solution V, the volumetric factor γ of dimension M⁻¹ is:

```
\gamma: M^{-1} = N_A V where N_A: mol^{-1} and V:L 
#X / \gamma: M = concentration of X molecules 
\gamma \cdot [X]: 1 = total number of X molecules (rounded to an integer).
```

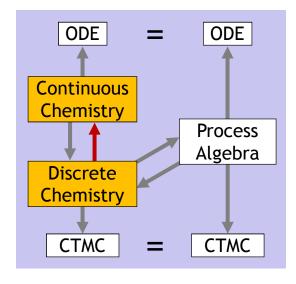
The Gillespie Conversion

Discrete Chemistry	Continuous Chemistry	$\gamma = N_A V$:M ⁻¹
initial quantities $\#A_0$	initial concentration [A] ₀	ns with [A] ₀ =#	A_0/γ
A ⊶ ^r A′	$A \rightarrow^k A'$	with $k = r$:s ⁻¹
A+B ⊶ A'+B'	$A+B \rightarrow^k A'+B'$	with $k = r\gamma$:M ⁻¹ s ⁻¹
A+A ⊶ ^r A'+A"	$A+A \rightarrow^k A'+A''$	with $k = r\gamma/2$:M ⁻¹ s ⁻¹

V = interaction volume N_A = Avogadro's number

Think
$$\gamma = 1$$
 i.e. $V = 1/N_A$

 $M = mol \cdot L^{-1}$ molarity (concentration)



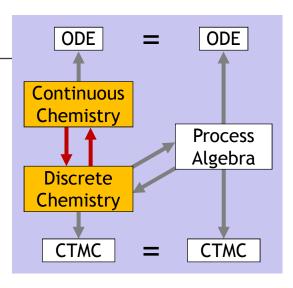
$Cont_{\gamma}$ and $Disc_{\gamma}$

4.2-3 Definition: Cont₇ and Disc₇

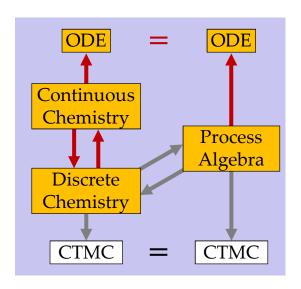
For a volumetric factor $\gamma:M^{-1}$, we define a translation $Cont_{\gamma}$ from a discrete chemical systems (C,P), with species X and initial molecule count $\#X_0 = \#X(P)$, to a continuous chemical systems (C,V) with initial concentration $[X]_0 = V_X$. The translation $Disc_{\gamma}$ is its inverse, up to a rounding error $\lceil \gamma[X]_0 \rceil$ in converting concentrations to molecule counts. Since γ is a global conversion constant, we later usually omit it as a subscript.

$Cont_{\gamma}(X \rightarrow^{r} P)$	$= X \rightarrow^k P$	with $k = r$,	r:s ⁻¹	k:s ⁻¹
$Cont_{\gamma}(X+Y \rightarrow^{\mathbf{r}} P)$	$= X+Y \rightarrow^k P$	with $k = r\gamma$	r:s ⁻¹	k:M ⁻¹ s ⁻¹
$Cont_{\gamma}(X+X \rightarrow^{\mathbf{r}} P)$	$= X+X \rightarrow^k P$	with $k = r\gamma/2$	r:s ⁻¹	k:M ⁻¹ s ⁻¹
$Cont_{\gamma}(\#X_0)$	$= [X]_0$	with $[X]_0 = \#X_0/\gamma$	$X_0:mol$	$[X]_0:M$
$Disc_{\gamma}(X \rightarrow^{k} P)$	$= X \rightarrow^{r} P$	with $r = k$,	k:s ⁻¹	r:s ⁻¹
$Disc_{\gamma}(X \to^{k} P)$ $Disc_{\gamma}(X+Y \to^{k} P)$	$= X \rightarrow^{r} P$ $= X+Y \rightarrow^{r} P$	with $r = k$, with $r = k/\gamma$	k:s ⁻¹ k:M ⁻¹ s ⁻¹	
**		·		r:s ⁻¹

 $Ch_{\gamma} := Cont_{\gamma} \circ Ch$



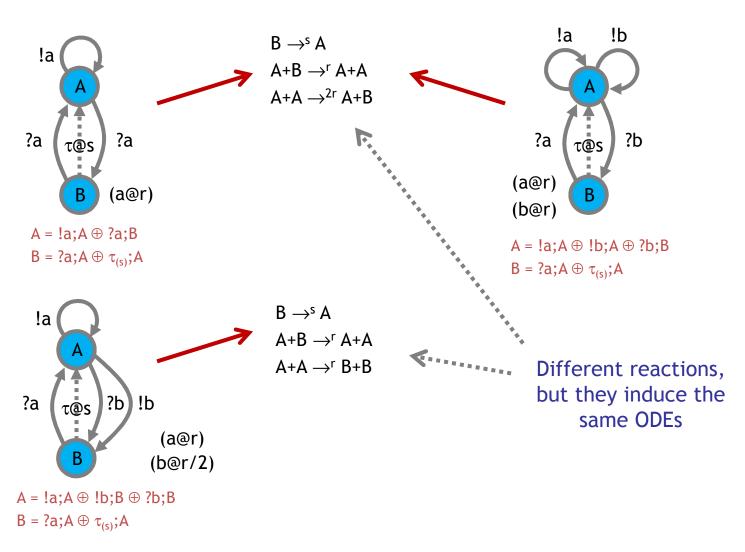
Continuous-State Semantics



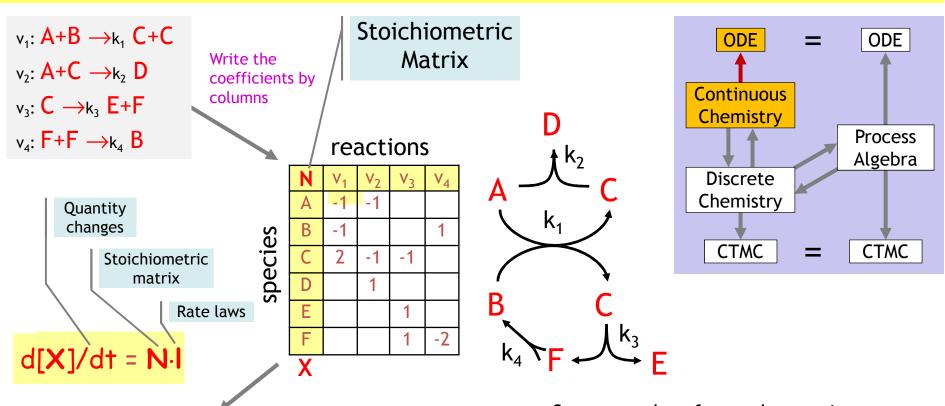
Same Semantics

Could chemistry itself be that semantics?

No: different sets of reactions can have the same behavior!



From Reactions to ODEs (Law of Mass Action)



$$d[A]/dt = -l_1 - l_2$$

$$d[B]/dt = -l_1 + l_4$$

$$d[C]/dt = 2l_1 - l_2 - l_3$$

$$d[D]/dt = l_2$$

$$d[E]/dt = l_3$$

$$d[F]/dt = l_3 - 2l_4$$

Read the concentration changes from the rows

E.g.
$$d[A]/dt = -k_1[A][B] - k_2[A][C]$$

Set a rate law for each reaction (Degradation/Hetero/Homeo)



X: chemical species

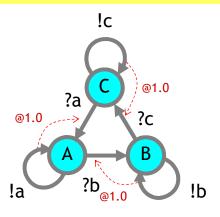
[-]: quantity of molecules

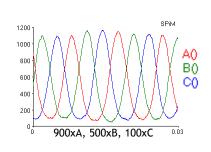
I: rate laws

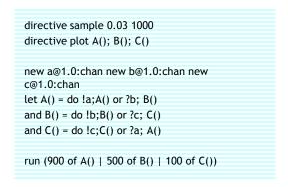
k: kinetic parameters

 ${\bf N}$: stoichiometric matrix

From Processes to ODEs via Chemistry!



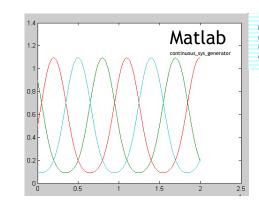




```
A = !a_{(s)}; A \oplus ?b_{(s)}; B
B = !b_{(s)}; B \oplus ?c_{(s)}; C
C = !c_{(s)}; C \oplus ?a_{(s)}; A
```

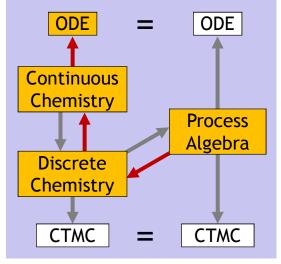
$$A+B \rightarrow^s B+B$$

 $B+C \rightarrow^s C+C$
 $C+A \rightarrow^s A+A$

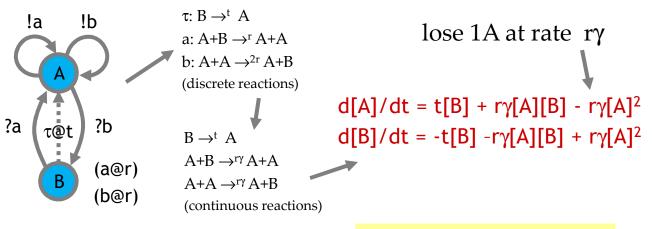


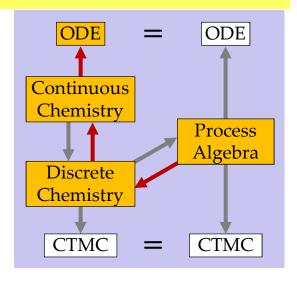
```
d[A]/dt = -s[A][B]+s[C][A]
d[B]/dt = -s[B][C]+s[A][B]
d[C]/dt = -s[C][A]+s[B][C]
```

$$(\gamma = 1)$$

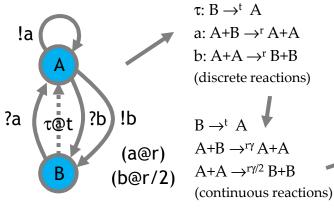


From Processes to ODEs via Chemistry!





Different chemistry but same ODEs, hence equivalent automata



lose 2A at rate
$$r\gamma/2$$

$$d[A]/dt = t[B] + r\gamma[A][B] - r\gamma[A]^{2}$$

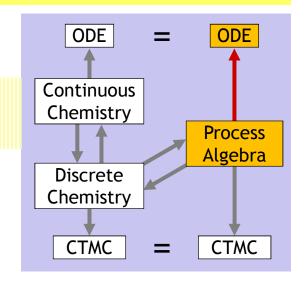
$$d[B]/dt = -t[B] - r\gamma[A][B] + r\gamma[A]^{2}$$

Processes Rate Equation

Process Rate Equation for Reagents E in volume γ

$$d[X]/dt = (\Sigma(Y \in E) Accr_{E}(Y,X) \cdot [Y]) - Depl_{E}(X) \cdot [X]$$
for all $X \in E$

"The change in process concentration (!!) for X at time t is:
the sum over all possible (kinds of) processes Y of:
the concentration at time t of Y
times the accretion from Y to X
minus the concentration at time t of X
times the depletion of X to some other Y"



$$Depl_{E}(X) =$$

$$\Sigma$$
(i: E.X.i= $\tau_{(r)}$;P) r +

$$\Sigma$$
(i: E.X.i=? $a_{(r)}$;P) r γ ·OutsOn_F(a) +

$$\Sigma$$
(i: E.X.i=!a_(r);P) r γ ·InsOn_E(a)

$$Accr_{E}(Y, X) =$$

$$\Sigma$$
(i: E.Y.i=t_(r);P) #X(P)·r +

$$\Sigma$$
(i: E.Y.i=? $a_{(r)}$;P) #X(P)·r γ ·OutsOn_E(a) +

$$\Sigma$$
(i: E.Y.i=! $a_{(r)}$;P) #X(P)·r γ ·InsOn_E(a)

InsOn_E(a) =
$$\Sigma$$
(Y \in E) #{Y.i | E.Y.i=?a_(r);P}·[Y]
OutsOn_E(a) = Σ (Y \in E) #{Y.i | E.Y.i=!a_(r);P}·[Y]

$$X = \tau_{(r)}; 0 \longrightarrow d[X]/dt = -r[X]$$

$$X = a_{(r)};0$$

$$d[X]/dt = -r\gamma[X][Y]$$

$$Y = !a_{(r)};0$$
 $d[Y]/dt = -r\gamma[X][Y]$

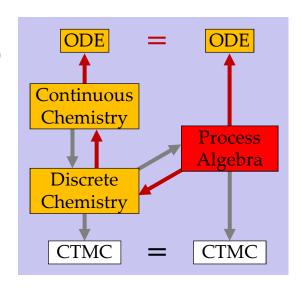
$$X = ?a_{(r)};0 \longrightarrow d[X]/dt = -2r\gamma[X]^2$$
 $\oplus !a_{(r)};0$

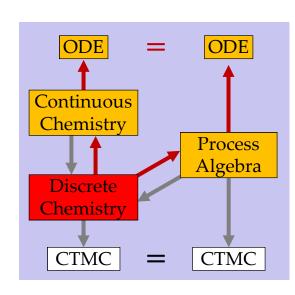
Continuous State Equivalence

Def: ≈ is equivalence of polynomials over the field of reals.

• Thm: $E \approx Cont(Ch(E))$

• Thm: Cont(C) \approx Pi(C)

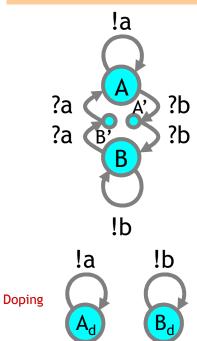




- For each E there is an E' \approx E that is detangled (E' = Pi(Ch(E)))
- For each E in automata form there is an an E' ≈ E that is detangled and in automata form (E' = Detangle(E)).

Exercise 2

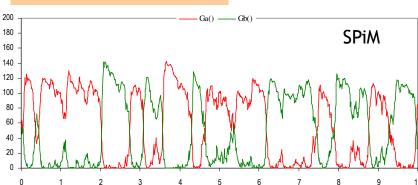
Q: What does this do?



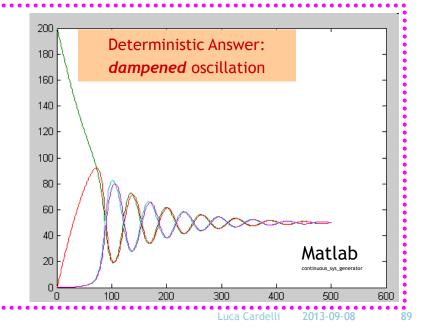
$$A = !a_{(r)}; A \oplus ?b; A'$$
 $A' = ?b; B$ $A_d = !a_{(r)}; A_d$ $B = !b_{(r)}; B \oplus ?a; B'$ $B' = ?a; A$ $B_d = !b_{(r)}; B_d$

Derive the ODEs from these "Hysteric Groupies" automata. Either by going through the chemical reactions and the Law of Mass Action (easier), or directly from the Process Rate Equation.

Stochastic Answer: robust quasi-oscillation



ODE predicts dampened oscillation, while the stochasic system keeps oscillating at max level.





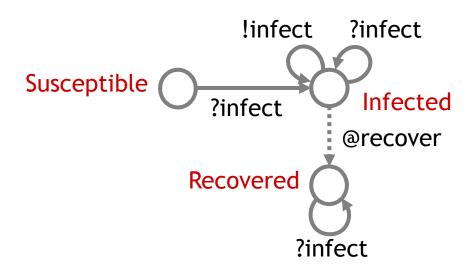
Epidemics

Non-Chemical Mass Action

Kermack, W. O. and McKendrick, A. G. "A Contribution to the Mathematical Theory of Epidemics." *Proc. Roy. Soc. Lond. A* 115, 700-721, 1927.

http://mathworld.wolfram.com/Kermack-McKendrickModel.html

Epidemics



Developing the Use of Process Algebra in the Derivation and Analysis of Mathematical Models of Infectious Disease

R. Norman and C. Shankland

Department of Computing Science and Mathematics, University of Stirling, UK. {ces,ran}@cs.stir.ac.uk

Abstract. We introduce a series of descriptions of disease spread using the process algebra WSCCS and compare the derived mean field equations with the traditional ordinary differential equation model. Even the preliminary work presented here brings to light interesting theoretical questions about the "best" way to defined the model.

directive sample 500.0 1000
directive plot Recovered(); Susceptible(); Infected()

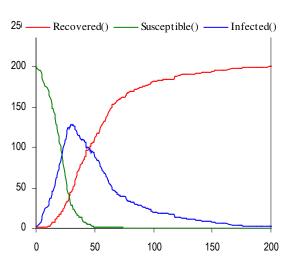
new infect @0.001:chan()
val recover = 0.03

let Recovered() =
 ?infect; Recovered()

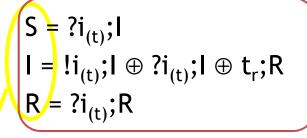
and Susceptible() =
 ?infect; Infected()

and Infected() =
 do !infect; Infected()
 or ?infect; Infected()
 or delay@recover; Recovered()

run (200 of Susceptible() | 2 of Infected())



ODEs



Differentiating

"useless" reactions

$$d[S]/dt = -t\gamma[S][I]$$

 $d[I]/dt = t\gamma[S][I]-r[I]$
 $d[R]/dt = r[I]$

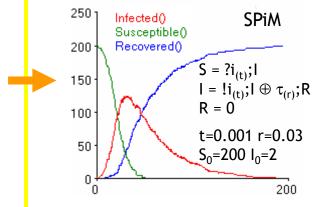
Automata produce the standard ODEs!

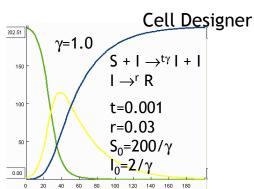
$$\frac{dS}{dt} = -aIS$$

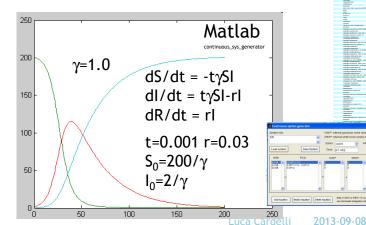
$$\frac{dI}{dt} = aIS - bI$$

$$\frac{dR}{dt} = bI$$

(the Kermack-McKendrick, or SIR model)!

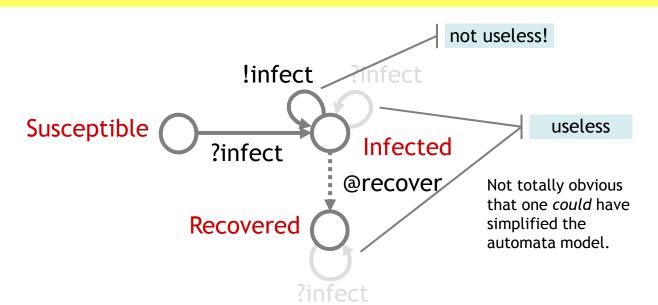


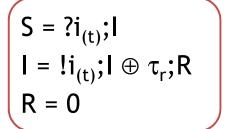




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Simplified Model



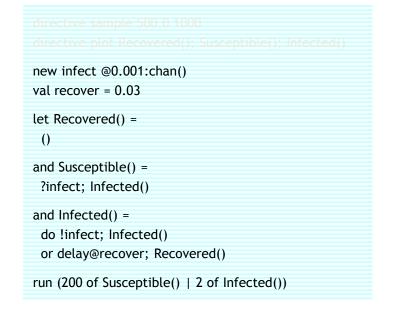


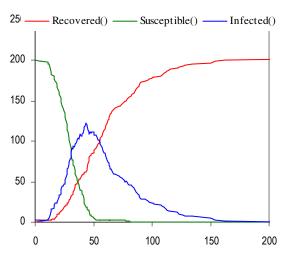
```
d[S]/dt = -t\gamma[S][I]

d[I]/dt = t\gamma[S][I]-r[I]

d[R]/dt = r[I]
```

Same ODE, hence equivalent automata models.

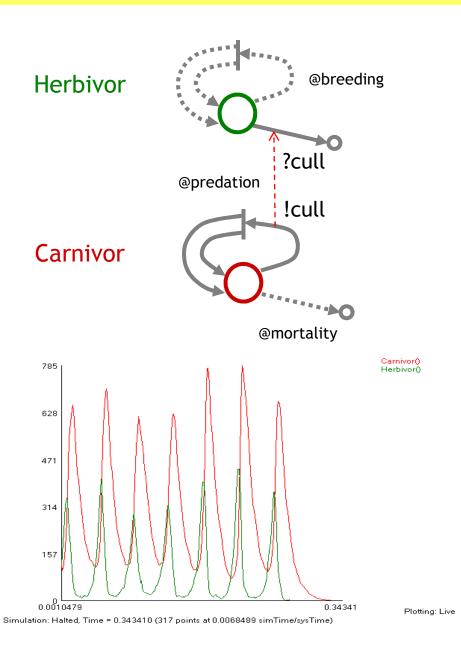


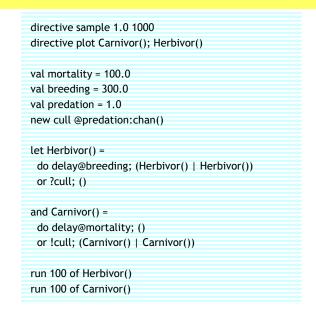


Lotka-Volterra

Unbounded Systems

Predator-Prey





An unbounded state system!

Lotka-Volterra in Matlab

$$H = \tau_{b}; (H|H) \oplus ?c_{(p)}; 0$$

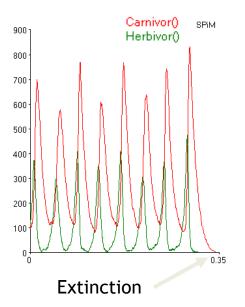
$$C = \tau_{m}; 0 \oplus !c_{(p)}; (C|C)$$

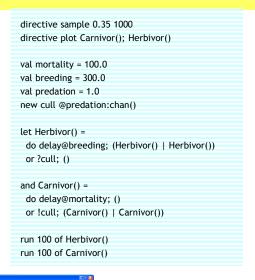
$$\#H_{0}, \#C_{0}$$

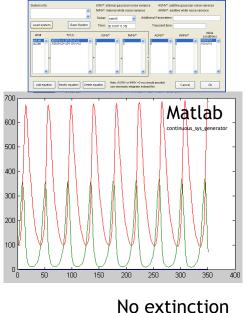
$$\begin{pmatrix} \mathsf{H} \to^{\mathsf{b}} \mathsf{H} + \mathsf{H} \\ \mathsf{C} \to^{\mathsf{m}} \mathsf{0} \\ \mathsf{H} + \mathsf{C} \to^{\mathsf{p}\gamma} \mathsf{C} + \mathsf{C} \\ [\mathsf{H}]_0 = \#\mathsf{H}_0/\gamma \\ [\mathsf{C}]_0 = \#\mathsf{C}_0/\gamma$$

 $d[H]/dt = b[H]-p\gamma[H][C]$ $d[C]/dt = -m[C]+p\gamma[H][C]$ $[H]_0 = \#H_0/\gamma$ $[C]_0 = \#C_0/\gamma$

```
m=100.0
b=300.0
p=1.0
\gamma=1.0
#H<sub>0</sub> = 100
#C<sub>0</sub> = 100
```

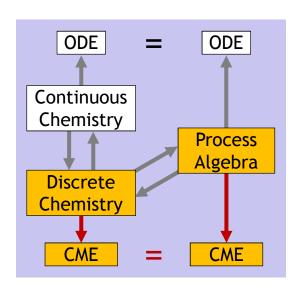






Which one is the "right prediction"?

Master Equation Semantics



Chemical Master Equation

Chemical Master Equation for a chemical system C

$$\partial pr(s,t)/\partial t = \sum_{i \in 1..M} a_i(s-v_i) \cdot pr(s-v_i,t) - a_i(s) \cdot pr(s,t)$$
 for all $s \in States(C)$

Reactions Propensity

"The change of probability at time t of a state is:

the sum over all possible (kinds of) reactions of:

the probability at time t of each state leading to this one
times the propensity of that reaction in that state
minus the probability at time t of the current state
times the propensity of each reaction in the current state"

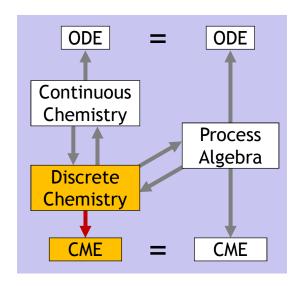
 $s \in 1..N \rightarrow Nat$ is a *state* of the system with N chemical species

pr(s,t) = Pr{ $\chi(t)$ =s | $\chi(0)$ =s₀} is the conditional probability of the system χ being in state s at time t given that it was in state s₀ at time 0.

There are 1...M chemical reactions.

v_i is the state change caused by reaction i (as a difference)

 $a_i(s) = c_i \cdot h_i(r)$ is the *propensity* of reaction i in state s, defined by a base reaction rate and a state-dependent count of the distinct combinations of reagents. (It depends on the kind of reactions.)



Process Algebra Master Equation

Process Master Equation for a system of reagents E

$$\partial \text{pr}(r,t)/\partial t = \sum_{i \in S} a_i(r-v_i)\cdot \text{pr}(r-v_i,t) - a_i(r)\cdot \text{pr}(r,t)$$
 for all $r \in \text{States}(E)$

Interactions Propensity

"The change of probability at time t of a state is:

the sum over all possible (kinds of) interactions of:

the probability at time t of each state leading to this one
times the propensity of that interaction in that state
minus the probability at time t of the current state
times the propensity of each interaction in the current state"

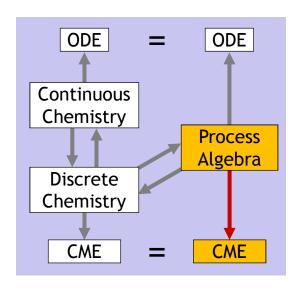
 $r \in species(E) \rightarrow Nat$ is a *state* of the system

pr(r,t) = Pr{ $\chi(t)$ =r | $\chi(0)$ =r₀} is the conditional probability of the system χ being in state r at time t given that it was in state r₀ at time 0.

 $\mathfrak S$ is the finite set of *possible interactions* arising from a set of reagents E. (All τ and all ?a/!a pairs in E)

v_i is the state change caused by interaction i (as a difference)

 $a_i(r) = r_i \cdot h_i(r)$ is the *propensity* of interaction i in state r, defined by a base rate of interaction and a state-dependent count of the distinct combinations of reagents. (It depends on the kind of interaction.)



... details

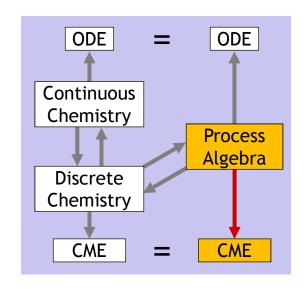
Process Master Equation for Reagents E

$$\partial pr(r,t)/\partial t = \sum_{i \in \mathcal{I}} a_i(r-v_i) \cdot pr(r-v_i,t) - a_i(r) \cdot pr(r,t)$$
 for all $r \in States(E)$

- $pr(p,t) = Pr\{S(t)=p \mid S(0)=p_0\}$ is the conditional probability of the system being in state p (a multiset of molecules) at time t given that it was in state p_0 at time 0.
- $$\label{eq:continuous} \begin{split} \mathfrak{I} &= \{\{X.i\} \ s.t. \ E.X.i = \tau_{(r)}; Q\} \ \cup \\ &\{\{X.i, \ Y.j\} \ s.t. \ E.X.i = ?n_{(r)}; Q \ and \ E.Y.j = !n_{(r)}; R\} \\ &\text{is the set of possible interactions in E} \end{split}$$

 v_i is the state change caused by an interaction $i \in \mathcal{I}$.

$$v_i = -X+Q$$
 if $i = \{X.i\}$ s.t. E.X. $i = \tau_{(r)}$; Q $v_i = -X-Y+Q+R$ if $i = \{X.i, Y.j\}$ s.t. E.X. $i = ?n_{(r)}$; Q and E.Y. $j = !n_{(r)}$; R



 a_i is the *propensity* of interaction i in state p. Here $p^{\#X}$ is the number of X in p.

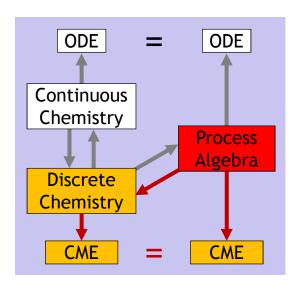
$$\begin{array}{ll} a_i(p) = r \cdot p^{\#X} & \text{if } i = \{X.i\} \ \textit{s.t.} \ E.X.i = \tau_{(r)}; Q \\ a_i(p) = r \cdot p^{\#X} \cdot p^{\#Y} & \text{if } i = \{X.i, \ Y.j\} \ \textit{s.t.} \ X \neq Y \ \text{and} \ E.X.i = ?a_{(r)}; Q \ \text{and} \ E.Y.j = !a_{(r)}; R \\ a_i(p) = r \cdot p^{\#X} \cdot (p^{\#X} - 1) & \text{if } i = \{X.i, \ X.j\} \ \textit{s.t.} \ E.X.i = ?a_{(r)}; Q \ \text{and} \ E.X.j = !a_{(r)}; R \end{array}$$

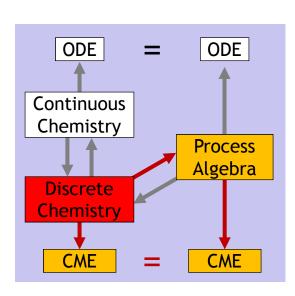
Equivalence of Master Equations

• Def: ≈ is equivalence of derived Master Equations (they are identical).

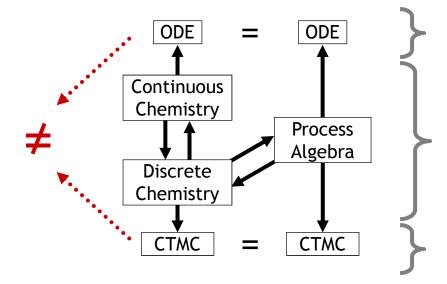
• Thm: $E \approx Ch(E)$

• Thm: $C \approx Pi(C)$





GMA ≠ CME



Semantics #1
Continuous state space

Syntax

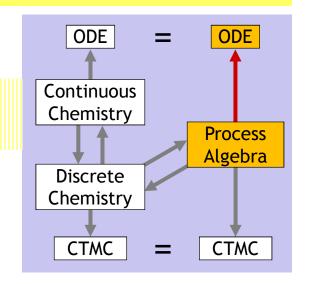
Semantics #2
Discrete state space

Processes to GMA Directly

Process Rate Equation for Reagents E in volume γ

$$d[X]/dt = (\Sigma(Y \in E) Accr_{E}(Y,X) \cdot [Y]) - Depl_{E}(X) \cdot [X]$$
for all $X \in E$

"The change in process concentration (!!) for X at time t is:
the sum over all possible (kinds of) processes Y of:
the concentration at time t of Y
times the accretion from Y to X
minus the concentration at time t of X
times the depletion of X to some other Y"



$$Depl_{E}(X) =$$

$$\Sigma$$
(i: E.X.i= τ _(r);P) r +

$$\Sigma$$
(i: E.X.i=? $a_{(r)}$;P) r γ ·OutsOn_E(a) +

$$\Sigma$$
(i: E.X.i=!a_(r);P) r γ ·InsOn_E(a)

$$Accr_{E}(Y, X) =$$

$$\Sigma$$
(i: E.Y.i=t_(r);P) #X(P)·r +

$$\Sigma$$
(i: E.Y.i=? $a_{(r)}$;P) #X(P)·r γ ·OutsOn_E(a) +

$$\Sigma$$
(i: E.Y.i=! $a_{(r)}$;P) #X(P)·r γ ·InsOn_E(a)

InsOn_E(a) =
$$\Sigma$$
(Y \in E) #{Y.i | E.Y.i=?a_(r);P}·[Y]
OutsOn_E(a) = Σ (Y \in E) #{Y.i | E.Y.i=!a_(r);P}·[Y]

$$X = \tau_{(r)}; 0 \longrightarrow d[X]/dt = -r[X]$$

$$X = a_{(r)};0$$
 $d[X]/dt = -r\gamma[X][Y]$

$$Y = !a_{(r)};0$$
 $d[Y]/dt = -r\gamma[X][Y]$

$$X = ?a_{(r)};0 \longrightarrow d[X]/dt = -2r\gamma[X]^2$$

$$\oplus !a_{(r)};0$$

Process Algebra Master Equation

Process Master Equation for a system of reagents E

$$\partial pr(r,t)/\partial t = \sum_{i \in \mathcal{I}} a_i(r-v_i) \cdot pr(r-v_i,t) - a_i(r) \cdot pr(r,t)$$
 for all $r \in States(E)$

Interactions Propensity

"The change of probability at time t of a state is:

the sum over all possible (kinds of) interactions of:

the probability at time t of each state leading to this one
times the propensity of that interaction in that state
minus the probability at time t of the current state
times the propensity of each interaction in the current state"

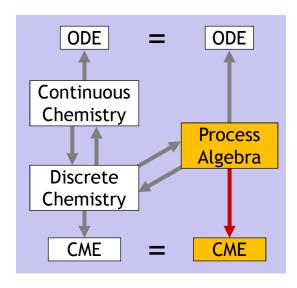
 $r \in species(E) \rightarrow Nat$ is a *state* of the system

pr(r,t) = Pr{ $\chi(t)$ =r | $\chi(0)$ =r₀} is the conditional probability of the system χ being in state r at time t given that it was in state r₀ at time 0.

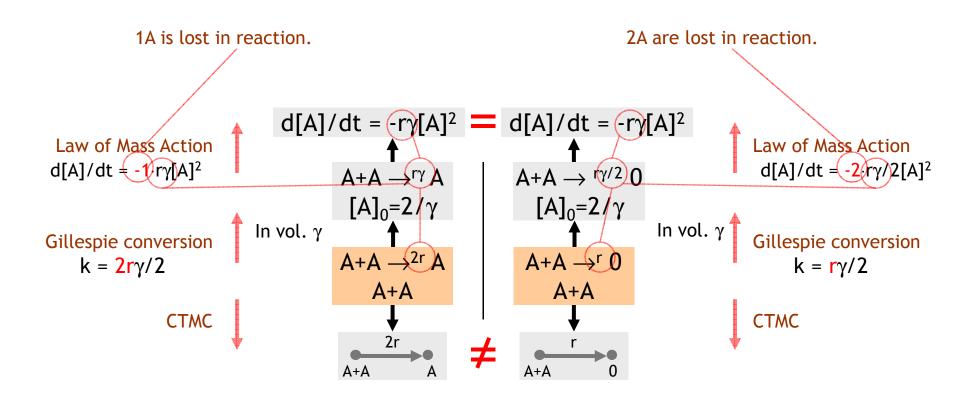
 $\mathfrak S$ is the finite set of *possible interactions* arising from a set of reagents E. (All τ and all ?a/!a pairs in E)

v_i is the state change caused by interaction i (as a difference)

 $a_i(r) = r_i \cdot h_i(r)$ is the *propensity* of interaction i in state r, defined by a base rate of interaction and a state-dependent count of the distinct combinations of reagents. (It depends on the kind of interaction.)

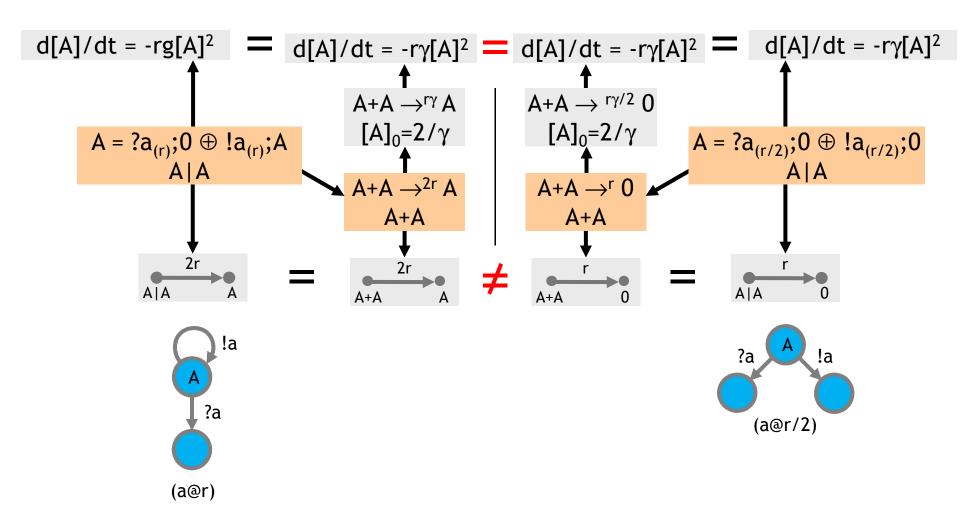


$A+A \rightarrow^{2r} A =? A+A \rightarrow^{r} 0$

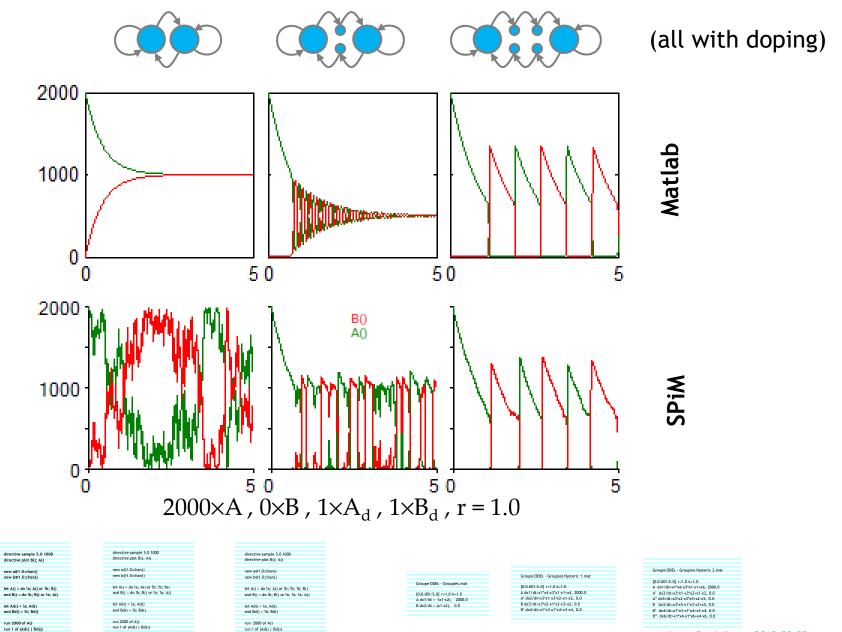


(For conservation of mass, consider instead $A+A \rightarrow^{2r} A+B$ vs. $A+A \rightarrow^{r} B+B$)

$A+A \rightarrow^{2r} A =? A+A \rightarrow^{r} 0$

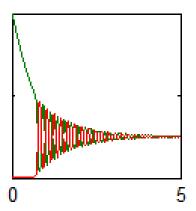


Continuous vs. Discrete Groupies

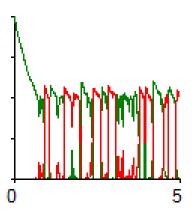


Scientific Predictions





After a while, all 4 states are almost equally occupied.

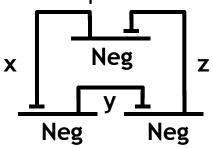


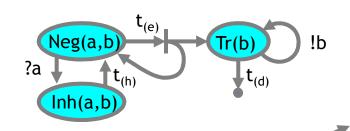
The 4 states are almost never equally occupied.

And Yet It Moves

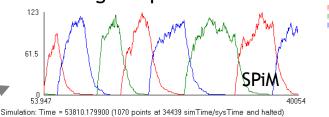
R.Blossey, L.Cardelli, A.Phillips: Compositionality, Stochasticity and Cooperativity in Dynamic Models of Gene Regulation (HFSP Journal)

The Repressilator





A fine stochastic oscillator over a wide range of parameters.



Paused

Parametric representation

```
\begin{split} &\text{Neg}(a,b) = ?a; \; \text{Inh}(a,b) \oplus \tau_e; \; (\text{Tr}(b) \mid \text{Neg}(a,b)) \\ &\text{Inh}(a,b) = \tau_h; \; \text{Neg}(a,b) \\ &\text{Tr}(b) = !b; \; \text{Tr}(b) \oplus \tau_g; \; 0 \\ &\text{Neg}(x_{(r)},y_{(r)}) \mid \text{Neg}(y_{(r)},z_{(r)}) \mid \text{Neg}(z_{(r)},x_{(r)}) \end{split}
```

```
Neg/x,y \rightarrow e Tr/y + Neg/x,y
```

$$Neg/y,z \rightarrow^e Tr/z + Neg/y,z$$

$$Neg/z,x \rightarrow^e Tr/x + Neg/z,x$$

$$Tr/x + Neg/x,y \rightarrow^r Tr/x + Inh/x,y$$

$$Tr/y + Neg/y,z \rightarrow^r Tr/y + Inh/y,z$$

$$Tr/z + Neg/z, x \rightarrow^r Tr/z + Inh/z, x$$

$$lnh/x,y \rightarrow^h Neg/x,y$$

Inh/y,z
$$\rightarrow$$
h Neg/y,z

$$lnh/z,x \rightarrow^h Neg/z,x$$

$$Tr/x \rightarrow g 0$$

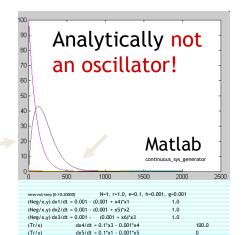
$$Tr/y \rightarrow g 0$$

$$Tr/z \rightarrow^g 0$$

$$Neg/x,y + Neg/y,z + Neg/z,x$$

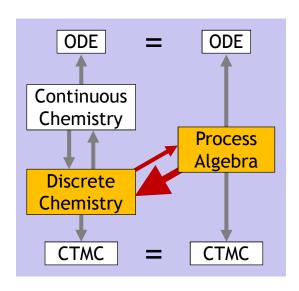
simplifying (N is the quantity of each of the 3 gates)

```
d[Neg/x,y]/dt = hN - (h+r[Tr/x])[Neg/x,y]
d[Neg/y,z]/dt = hN - (h+r[Tr/y])[Neg/y,z]
d[Neg/z,x]/dt = hN - (h+r[Tr/z])[Neg/z,x]
d[Tr/x]/dt = e[Neg/z,x] - g[Tr/x]
d[Tr/y]/dt = e[Neg/x,y] - g[Tr/y]
d[Tr/z]/dt = e[Neg/y,z] - g[Tr/z]
```



dx6/dt = 0.1*x2 - 0.001*x6

Model Compactness



n² Scaling Problems

StoichiometricMatrix($Ch(E_3)$)

- E_n has 2n variables (nodes) and 2n terms (arcs).

 $Ch(E_3)$

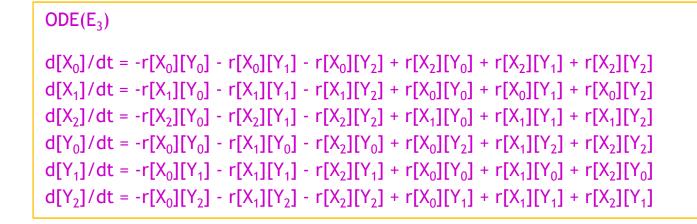
- $Ch(E_n)$ has 2n species and n^2 reactions.

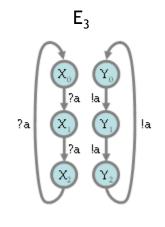
- The stoichiometric matrix has size $2n \cdot n^2 = 2n^3$.
- The ODEs have 2n variables and $2n(n+n) = 4n^2$ terms (number of variables times number of accretions plus depletions when sums are distributed)

E_3 $X_0 = ?a_{(r)}; X_1$ $X_1 = ?a_{(r)}; X_2$ $X_2 = ?a_{(r)}; X_0$ $Y_0 = !a_{(r)}; Y_1$ $Y_1 = !a_{(r)}; Y_2$ $Y_2 = !a_{(r)}; Y_0$

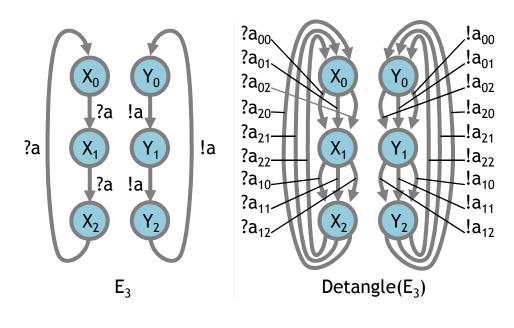
$\begin{array}{c} a_{00} \colon X_0 + Y_0 \to^r X_1 + Y_1 \\ a_{01} \colon X_0 + Y_1 \to^r X_1 + Y_2 \\ a_{02} \colon X_0 + Y_2 \to^r X_1 + Y_0 \\ a_{10} \colon X_1 + Y_0 \to^r X_2 + Y_1 \\ a_{11} \colon X_1 + Y_1 \to^r X_2 + Y_2 \\ a_{12} \colon X_1 + Y_2 \to^r X_2 + Y_0 \\ a_{20} \colon X_2 + Y_0 \to^r X_0 + Y_1 \\ a_{21} \colon X_2 + Y_1 \to^r X_0 + Y_2 \\ a_{22} \colon X_2 + Y_2 \to^r X_0 + Y_0 \end{array}$

	a ₀₀	a ₀₁	a ₀₂	a ₁₀	a ₁₁	a ₁₂	a ₂₀	a ₂₁	a ₂₂
Xo	-1							+1	
X_1			+1	-1	-1	-1			
X ₂				+1	+1	+1	-1	-1	-1
y ₀	-1		+1	-1		+1	-1		+1
y ₁	+1	-1		+1	-1		+1	-1	
y ₂		+1	-1		+1	-1		+1	-1





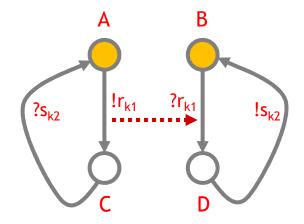
Entangled vs detangled



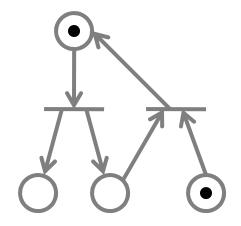
(closely related to $Pi(Ch(E_3))$)

Model Maintenance

- Biology (unlike much of chemistry) is combinatorial
 - Biochemical systems have many regular repeated components
 - Components interact and combine in complex combinatorial ways
 - Components have local state
 - A biochemical system is vastly more compact that its potential state space
- One may have to expand the state space during analysis, but must not do it during description
- There is a good way:
 - Describe biochemical systems compositionally
 - Each component with its own state and interactions
 - o ... as Nature intended...



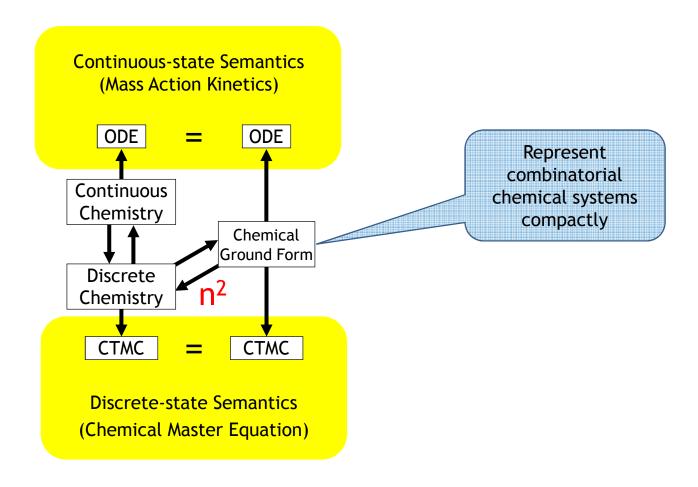
Or



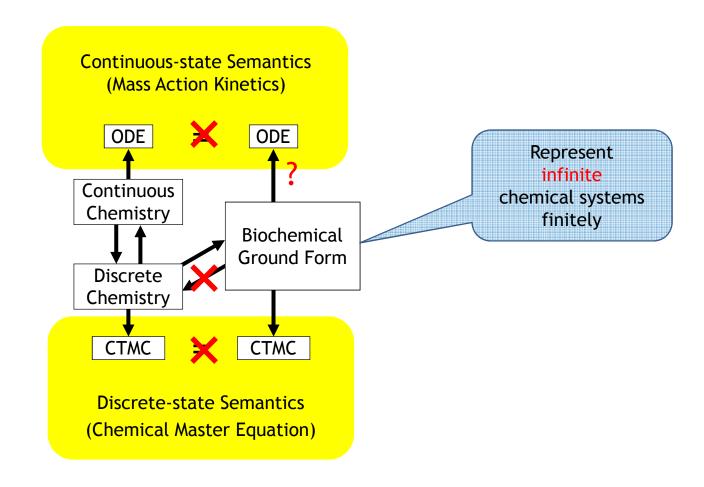
Or ...

Chemistry and Beyond

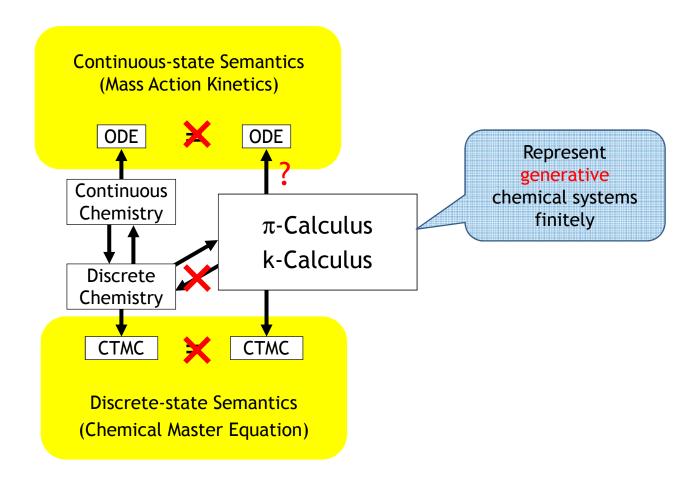
Process Algebra is 'Bigger' than Chemistry



Process Algebra is 'Bigger' than Chemistry



Process Algebra is 'Bigger' than Chemistry



Conclusions

Conclusions

Process Algebra

- An extension of automata theory to populations of interacting automata
- o Modeling the behavior of individuals in an arbitrary environment
- Compositionality (combining models by juxtaposition)

Connections between modeling approaches

- Connecting the discrete/concurrent/stochastic/molecular approach
- o to the continuous/sequential/deterministic/population approach

Connecting syntax with semantics

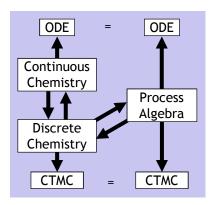
- Syntax = model presentation (equations/programs/diagrams/blobs etc.)
- Semantics = state space (generated by the syntax)

Ultimately, connections between analysis techniques

- We need (and sometimes have) good semantic techniques to analyze state spaces (e.g. calculus, but also increasingly modelchecking)
- But we need equally good syntactic techniques to structure complex models (e.g. compositionality) and analyze them (e.g. process algebra)

A bright future for Computer Science and Logic in modern Biology

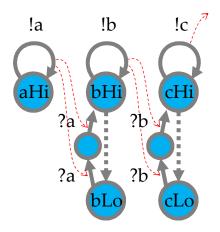
 Biology needs good analysis techniques for discrete systems analysis (modal logics, modelchecking, causality analysis, abstract interpretation, ...)



Exercise 1a Solution

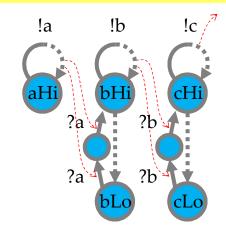
$$aHi = !a_{(1.0)}; aHi \\ bLo = ?a_{(1.0)}; bMd \\ bMd = ?a_{(1.0)}; bHi \\ bHi = !b_{(1.0)}; bHi \oplus \tau_{(1.0)}; bLo \\ cLo = ?b_{(1.0)}; cMd \\ cMd = ?b_{(1.0)}; cHi \\ cHi = !c_{(1.0)}; cHi \oplus \tau_{(1.0)}; cLo$$

aHi + bLo
$$\rightarrow^{1.0}$$
 aHi + bMd
aHi + bMd $\rightarrow^{1.0}$ aHi + bHi
bHi $\rightarrow^{1.0}$ bLo
bHi + cLo $\rightarrow^{1.0}$ bHi + cMd
bHi + cMd $\rightarrow^{1.0}$ bHi + cHi
cHi $\rightarrow^{1.0}$ cLo



Exercise 1b Solution

aHi =
$$!a_{(1.0)}$$
; aR
aR = $\tau_{(1.0)}$; aHi
bLo = $?a_{(1.0)}$; bMd
bMd = $?a_{(1.0)}$; bHi
bHi = $!b_{(1.0)}$; bR $\oplus \tau_{(1.0)}$; bLo
bR = $\tau_{(1.0)}$; bHi
cLo = $?b_{(1.0)}$; cMd
cMd = $?b_{(1.0)}$; cHi
cHi = $!c_{(1.0)}$; cR $\oplus \tau_{(1.0)}$; cLo
cR = $\tau_{(1.0)}$; cHi



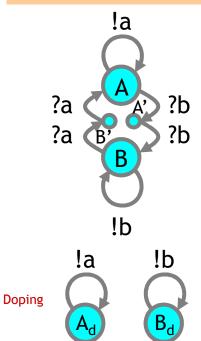
aHi + bLo
$$\rightarrow^{1.0}$$
 aR + bMd
aHi + bMd $\rightarrow^{1.0}$ aR + bHi
aR $\rightarrow^{1.0}$ aHi
bHi $\rightarrow^{1.0}$ bLo
bHi + cLo $\rightarrow^{1.0}$ bR + cMd
bHi + cMd $\rightarrow^{1.0}$ bR + cHi
bR $\rightarrow^{1.0}$ bHi
cHi $\rightarrow^{1.0}$ cLo

Note: no reaction from cHi to cR etc. because there is nothing (here) to interact with c.

The chemical system is incomplete (it does not say how cHi would behave in a bigger system), while the automata already specify what would happen (if we remove the red bits above we obtain the same reactions).

Exercise 2 Solution

Q: What does this do?



$$A = !a_{(r)}; A \oplus ?b; A' \quad A' = ?b; B \qquad A_d = !a_{(r)}; A_d$$

 $B = !b_{(r)}; B \oplus ?a; B' \quad B' = ?a; A \qquad B_d = !b_{(r)}; B_d$

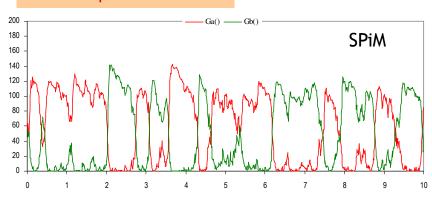
d[A]/dt = r[A][B']-r[B][A]-r[A][B_d]+r[B'][A_d]
d[A']/dt = r[B][A]-r[B][A']+r[A][B_d]-r[A'][B_d]
d[B]/dt = r[B][A']-r[A][B]-r[B][A_d]+r[A'][B_d]
d[B']/dt = r[A][B]-r[A][B']+r[B][A_d]-r[B'][A_d]

 $d[A_d]/dt = 0$ $d[B_d]/dt = 0$

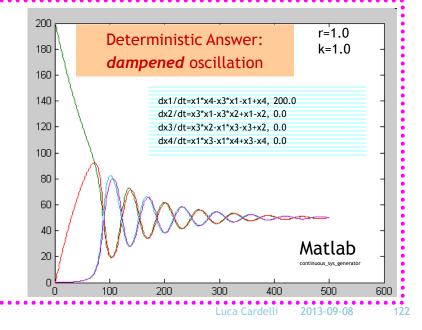
d[A]/dt = r[A][B']-r[B][A]-rk[A]+rk[B'] d[A']/dt = r[B][A]-r[B][A']+rk[A]-rk[A'] d[B]/dt = r[B][A']-r[A][B]-rk[B]+rk[A'] d[B']/dt = r[A][B]-r[A][B']+rk[B]-rk[B']

 $[A_d]$, $[B_d]$ are constant; assume them both = k

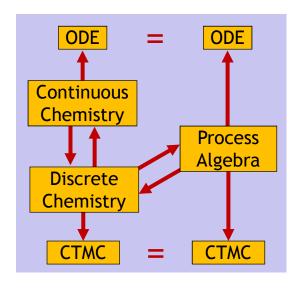
Stochastic Answer: robust quasi-oscillation



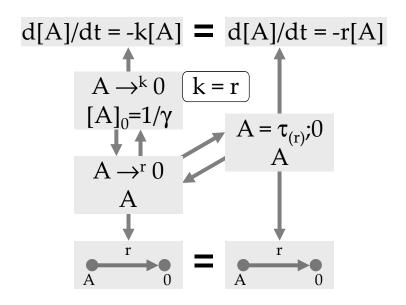
ODE predicts dampened oscillation, while the stochasic system keeps oscillating at max level.



Summary



Unary Reactions



O----- Unary Reaction

Hetero Reactions

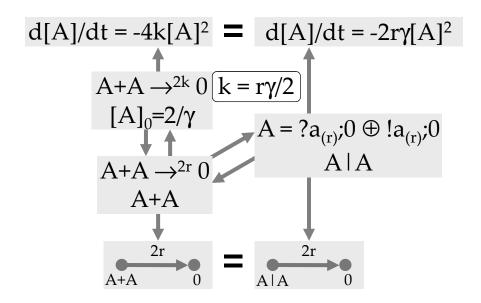
$$d[A]/dt=d[B]/dt=-k[A][B] = d[A]/dt=d[B]/dt=-r\gamma[A][B]$$

$$A+B \rightarrow^{k} 0 \quad k = r\gamma$$

$$[A]_{0}=[B]_{0}=1/\gamma$$

$$A+B \rightarrow^{r} 0$$

Homeo Reactions

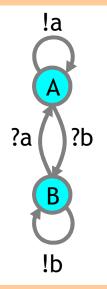


Homeo Reaction

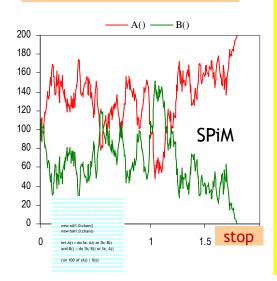
Groupies ODEs

Groupies ODE

Q: What does this do?



Stochastic Answer: **bistable** system



Same rate r

$$A = !a_{(r)};A \oplus ?b_{(r)};B$$

$$B = !b_{(r)};B \oplus ?a_{(r)};A$$

$$A+B \rightarrow^r A+A$$

 $B+A \rightarrow^r B+B$

$$d[A]/dt = r[A][B]-r[B][A]$$

 $d[B]/dt = r[B][A]-r[A][B]$

$$d[A]/dt = 0$$

 $d[B]/dt = 0$

Deterministic Answer: constant system

ODE predicts stability [A] = 0 for any value of [A], while the stochastic system is stable only when [A] is either 0 or Max.

Different rates r,s

$$A = !a_{(r)}; A \oplus ?b_{(s)}; B$$

$$B = !b_{(s)}; B \oplus ?a_{(r)}; A$$

$$\begin{pmatrix}
A+B \rightarrow^r A+A \\
B+A \rightarrow^s B+B
\end{pmatrix}$$

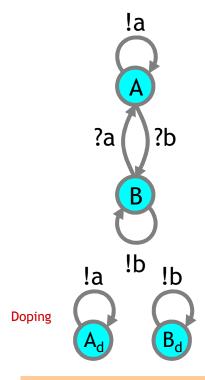
Deterministic Answer: *monostable* systems

ODE predicts A will always "win", if r>s, and B will always loose. The stochastic system can have either outcome (with different probabilities).

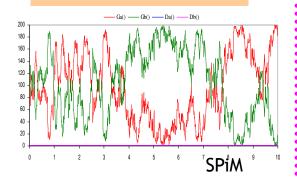


Doped Groupies ODEs

Q: What does this do?



Stochastic Answer: bounded random walk



$$A = !a_{(r)}; A \oplus ?b_{(r)}; B$$

$$B = !b_{(r)}; B \oplus ?a_{(r)}; A$$

$$A_d = !a_{(r)}; A_d$$

$$B_d = !b_{(r)}; B_d$$

$$A+B \rightarrow^r A+A$$
 $B+A \rightarrow^r B+B$ B

$$A+B_d \rightarrow^r B+B_d$$

$$B+A_d \rightarrow^r A+A_d$$

$$d[A]/dt = r[A][B]-r[B][A]-r[A][B_d]+r[B][A_d]$$

 $d[B]/dt = r[B][A]-r[A][B]-r[B][A_d]+r[A][B_d]$

$$d[A_d]/dt = 0$$

$$d[B_d]/dt = 0$$

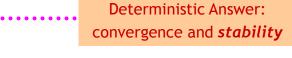
$$d[A]/dt = -rk([A]-[B])$$

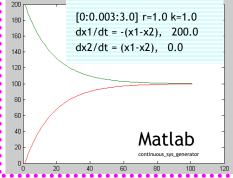
 $d[B]/dt = rk([A]-[B])$

At
$$[B]=0$$
: $d[A]/dt=-rk[A]$,

 $[A_d]$, $[B_d]$ are constant; assume them both = kd[B]/dt=rk[A]

At $[A] \approx [B]$: $d[A]/dt = d[B]/dt \approx 0$ At [A]=[B]: d[A]/dt=d[B]/dt=0



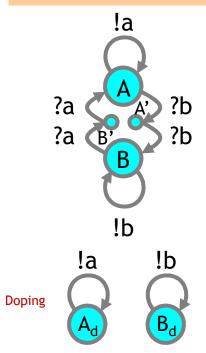


ODE predicts converging stable equilibrium at [A]=[B] instead of the total chaos observed in the stochastic system!

For k=0 (no dope), predicts deadlock d[A]/dt=d[B]/dt=0 but at any value of [A], which is definitely not true in the stochastic system.

Hysteric Groupies ODEs

Q: What does this do?



new a81.0:chan()
new b81.0:chan()
let Ga() = do Is; Ga() or ib; 7b; Gb()
and Gb() = do Is; Gb() or ib; 7b; Ga()
let CB() = 1c; Db()
and Db() = 1b; Db()
run 100 of (Ga() | Cb())
run 100 of (Ga() | Db())

$$A = !a_{(r)}; A \oplus ?b; A' \quad A' = ?b; B$$

 $B = !b_{(r)}; B \oplus ?a; B' \quad B' = ?a; A$

$$A_d = !a_{(r)}; A_d$$

$$B_d = !b_{(r)}; B_d$$

$$A+B \rightarrow^r A+B'$$
 $A+B' \rightarrow^r A+A$
 $B+A \rightarrow^r B+A'$ $B+A' \rightarrow^r B+B$

$$\begin{vmatrix} A+B_d \rightarrow^r & A'+B_d & A'+B_d \rightarrow^r & B+B_d \\ B+A_d \rightarrow^r & B'+A_d & B'+A_d \rightarrow^r & A+A_d \end{vmatrix}$$

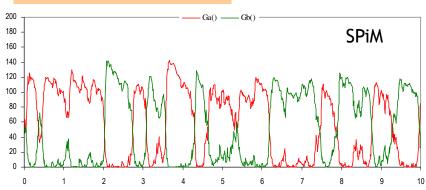
d[A]/dt = r[A][B']-r[B][A]-r[A][B_d]+r[B'][A_d]
d[A']/dt = r[B][A]-r[B][A']+r[A][B_d]-r[A'][B_d]
d[B]/dt = r[B][A']-r[A][B]-r[B][A_d]+r[A'][B_d]
d[B']/dt = r[A][B]-r[A][B']+r[B][A_d]-r[B'][A_d]

$$d[A_d]/dt = 0$$
$$d[B_d]/dt = 0$$

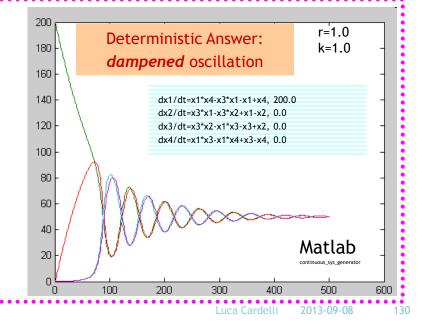
d[A]/dt = r[A][B']-r[B][A]-rk[A]+rk[B']
d[A']/dt = r[B][A]-r[B][A']+rk[A]-rk[A']
d[B]/dt = r[B][A']-r[A][B]-rk[B]+rk[A']
d[B']/dt = r[A][B]-r[A][B']+rk[B]-rk[B']

 $[A_d]$, $[B_d]$ are constant; assume them both = k

Stochastic Answer: robust quasi-oscillation

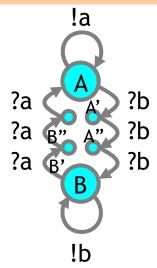


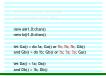
ODE predicts dampened oscillation, while the stochasic system keeps oscillating at max level.



Hysteric Groupies ODEs

Q: What does this do?





$$A = !a_{(r)}; A \oplus ?b; A' \quad A' = ?b; A'' \quad A'' = ?b; B$$

 $B = !b_{(r)}; B \oplus ?a; B' \quad B' = ?a; B'' \quad B'' = ?a; A$

$$A_d = !a_{(r)}; A_d$$

 $B_d = !b_{(r)}; B_d$

$$A+B \xrightarrow{\text{restored}(a) \mid \text{lock}(b) \mid \text{restored}(a) \mid \text{lock}(b) \mid \text{restored}(a) \mid \text{rest$$

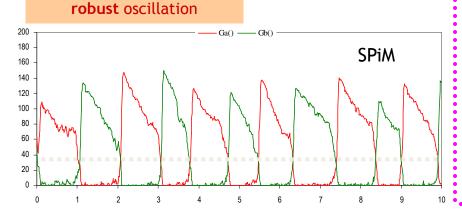
d[A]/dt = r[A][B"]-r[B][A]-r[A][B_d]+r[B"][A_d]
d[A']/dt = r[B][A]-r[B][A']+r[A][B_d]-r[A'][B_d]
d[A"]/dt = r[B][A']-r[B][A"]+r[A'][B_d]-r[A"][B_d]
d[B]/dt = r[B][A"]-r[A][B]-r[B][A_d]+r[A"][B_d]
d[B']/dt = r[A][B]-r[A][B']+r[B][A_d]-r[B'][A_d]
d[B"]/dt = r[A][B']-r[A][B"]+r[B'][A_d]-r[B"][A_d]

 $d[A_d]/dt = 0$ $d[B_d]/dt = 0$

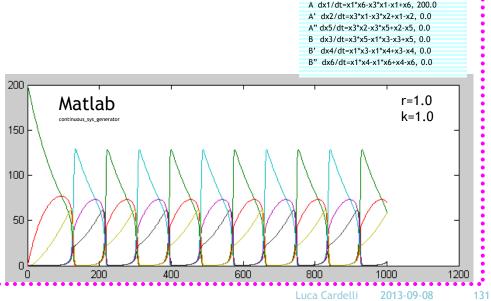
d[A]/dt = r[A][B"]-r[B][A]-rk[A]+rk[B"]
d[A']/dt = r[B][A]-r[B][A']+rk[A]-rk[A']
d[A"]/dt = r[B][A']-r[B][A"]+rk[A']-rk[A"]
d[B]/dt = r[B][A"]-r[A][B]-rk[B]+rk[A"]
d[B']/dt = r[A][B]-r[A][B']+rk[B]-rk[B']
d[B"]/dt = r[A][B']-r[A][B"]+rk[B']-rk[B"]

la !b Doping

Stochastic Answer:



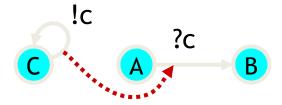
 $[A_d]$, $[B_d]$ are constant; assume them both = k



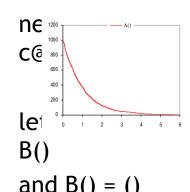
--Laws by ODEs

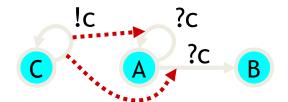
Choice Law by ODEs

Idle Interaction Law by ODEs



directive sample 6.0 1000 directive plot A()

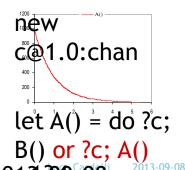




It may seem like A should decrease half as fast, but NO! Two ways to explain:

- -State A is *memoryless* of any past idling.
- Activity on c is double

directive sample 6.0 1000 directive plot A()



Equiconfluence Law by ODEs

$$t_{\lambda};B \mid t_{\lambda};B = t_{2\lambda};(B \mid t_{\lambda};B)$$

A t_{λ}

B

 t_{λ}

E

 $t_{2\lambda}$

G

B

$$A = t_{\lambda}; B$$

$$C = t_{\lambda}; B$$

$$A \mid C$$

$$E = t_{2\lambda}; (B \mid G)$$

 $G = t_{\lambda}; B$

$$\begin{bmatrix} E & \mathbb{R}^{2\lambda} & B + G \\ G & \mathbb{R}^{\lambda}B \end{bmatrix}$$

$$[E]' = -2\lambda[E]$$

$$[G]' = 2\lambda[E] - \lambda[G]$$

$$[B]' = 2\lambda[E] + \lambda[G]$$

$$[A']^{\cdot} = -\lambda[A']$$

$$[C']^{\cdot} = -\lambda[C']$$

$$[B]^{\cdot} = \lambda[A'] + \lambda[C']$$

let A' = C' = E+G/2

$$\lambda[A'] + \lambda[C']$$

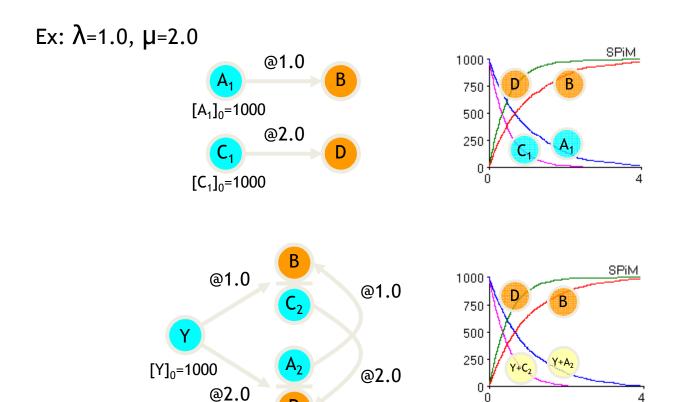
= $\lambda[E+G/2] + \lambda[E+G/2]$
= $2\lambda[E] + \lambda[G] = [B]$
[A']' = [E+G/2]' = [E]'+[G]'/2
= $-2\lambda[E] + (2\lambda[E] - \lambda[G])/2$
= $-\lambda[E] - \lambda[G]/2$
= $-\lambda[E+G/2] = -\lambda[A']$

[B] has equal time evolutions on the two sides provided that [A]=[E+G/2] and [C]=[E+G/2]. This imposes the constraint, in particular, that $[A]_0$ =[E+G/2] $_0$ and $[C]_0$ =[E+G/2] $_0$ (at time zero). The initial conditions of the right hand system specify that $[G/2]_0$ =0 (since only E is present). Therefore, we obtain that $[A]_0$ =[C] $_0$ =[E] $_0$.

2013-09-08

Stochastic Interleaving

$$t_{\lambda};B \mid t_{\mu};D = t_{\lambda};(B \mid t_{\mu};D) \text{ Å } t_{\mu};(t_{\lambda};B \mid D)$$



Amazingly, the B's and the D's from the two branches sum up to exponential distributions

```
directive sample 4.0 10000
directive plot A(); B(); C(); D()
let A() = delay@1.0; B()
and B() = ()
let C() = delay@2.0; D()
and D() = ()
run 1000 of (A() | C())
```

```
directive sample 4.0 10000
directive plot

?YA; B(); ?YC; D(); Y(); A(); C()
new YA@1.0:chan new YC@1.0:chan

let A() = do delay@1.0; B() or ?YA
and B() = ()

let C() = do delay@2.0; D() or ?YC
and D() = ()

let Y() =
   do delay@1.0; (B() | C())
   or delay@2.0; (A() | D())
   or ?YA or ?YC

run 1000 of Y()
```

Stochastic Interleaving Law by ODEs

$$t_{\lambda};B \mid t_{\mu};D = t_{\lambda};(B \mid t_{\mu};D) \text{ Å } t_{\mu};(t_{\lambda};B \mid D)$$

$$A_1 = t_{\lambda}; B$$

$$C_1 = t_{\mu}; D$$

$$A_1 \mid C_1$$

$$\begin{pmatrix}
A_1 & \mathbb{R}^{\lambda} & B \\
C_1 & \mathbb{R}^{\mu} & D \\
A_1 & + & C_1
\end{pmatrix}$$

$$\begin{bmatrix} [A^1] \\ [A^1] \end{bmatrix} = h[C^1]$$

$$\begin{bmatrix} [A^1] \\ [A^1] \end{bmatrix} = -\lambda[A^1]$$

$$Y = t_{\lambda}; (B \mid C_{2}) \text{ Å } t_{\mu}; (A_{2} \mid D)$$

$$C_{2} = t_{\mu}; D$$

$$A_{2} = t_{\lambda}; B$$

$$Y$$

$$\begin{pmatrix} Y & \mathbb{R}^{\lambda} & \mathbb{B} + \mathbb{C}_{2} \\ Y & \mathbb{R}^{\mu} & \mathbb{A}_{2} + \mathbb{D} \\ \mathbb{C}_{2} & \mathbb{R}^{\mu} & \mathbb{D} \\ \mathbb{A}_{2} & \mathbb{R}^{\lambda} & \mathbb{B} \\ Y \end{pmatrix} \longrightarrow \begin{pmatrix} [Y]^{\cdot} = -\lambda[Y] - \mu[Y] \\ [A_{2}]^{\cdot} = \mu[Y] - \lambda[A_{2}] \\ [B]^{\cdot} = \lambda[Y] + \lambda[A_{2}] \\ [C_{2}]^{\cdot} = \lambda[Y] - \mu[C_{2}] \\ [D]^{\cdot} = \mu[Y] + \mu[C_{2}]$$

$$\begin{cases} [Y+A_2]' = -\lambda[Y+A_2] \\ [B]' = \lambda[Y+A_2] \\ [Y+C_2]' = -\mu[Y+C_2] \\ [D]' = \mu[Y+C_2] \end{cases}$$

Want to show that B and D on both sides have the "same behavior" (equal quantities of B and D produced at all times)

$$[Y+A_2]' = [Y]'+[A_2]'$$

$$= -\lambda[Y]-\mu[Y]+\mu[Y]-\lambda[A_2]$$

$$= -\lambda[Y]-\lambda[A_2]$$

$$= -\lambda[Y+A_2] \qquad [Y+A_2] \text{ decays exponentially!}$$

[B] and [D] have equal time evolutions on the two sides provided that $[A_1]=[Y+A_2]$ and $[C_1]=[Y+C_2]$. This imposes the constraint, in particular, that $[A_1]_0=[Y+A_2]_0$ and $[C_1]_0=[Y+C_2]_0$ (at time zero). The initial conditions of the right hand system specify that $[A_2]_0=[C_2]_0=0$ (since only Y is present). Therefore, we obtain that $[A_1]_0=[C_1]_0=[Y]_0$.

So, for example, if we run a stochastic simulation of the left hand side with 1000*A1 and 1000*C1, we obtain the same curves for B and D than a stochastic simulation of the right hand side with 1000*Y.

2013-09-08